A geometry-aware compression of 3D mesh texture with random access

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Abstract—A 3D mesh object is usually represented as a combination of several entities including geometrical information (i.e., the triangles and their position in space) and a texture atlas/map (i.e., a giant 2D image containing all the texture information that is mapped to the 3D object at the rendering stage). This atlas is usually compressed using a conventional 2D image coder, thus without taking into account the geometrical information. Moreover, the whole image is usually decoded even though only a subpart of the mesh is observed by a user. In this paper, we propose a novel approach to compress a texture atlas of a 3D model that enables random access during decoding, and nevertheless takes into account the correlation driven by the geometrical information. The experimental results demonstrate the benefits of the proposed coder.

Index Terms—Random access, 3D mesh, texture map/atlas

I. INTRODUCTION

As shown in Fig. 1, a 3D model or object is usually rendered thanks to the combination of geometrical and textural information. This can be complemented by other types of information depicting for example the light reflection or granularity of the surface. Geometrical information is usually represented as a triangle mesh, defined by a list of triangles positioned in 3D space, and connected by their common edges. Efficient ways of compressing these meshes have been investigated in [1].

The textural information generally consists of a giant image, called texture map or atlas, describing the whole texture information that is mapped onto the 3D shape at the rendering stage. The correspondence between the texture and the triangles is also described in the mesh file. Texture images are usually coded as still images, with for instance JPEG [2] or even HEVC Intra [3]. This is however inefficient for the two following reasons.

First, a texture image is usually defined as an atlas of texture patches, with a flat/smooth background (see Fig. 1(b)). Even if optimizations of the patching arrangement have been proposed to lower the background size [4], [5], [6], [7], the projection of the texture content onto a 2D image introduces unavoidable artificial edges, and this leads to a rate overhead, when the encoding is performed with a regular image coder. Second, the atlas encoder inherits its property from conventional coder, i.e., the whole image has to be decoded. Nevertheless, this is not desirable since only a subpart of the 3D model, and therefore part of the atlas, is usually observed at a given time during the rendering. It would be much more efficient in terms of rate or memory load if only the needed part of the model could be decoded, as pointed in [8]. The question of random accessibility has been tackled for triangle meshes in [9]. In this work, we consider the problem of random accessibility for the texture information associated to such meshes.

A geometry-aware atlas coder has been proposed in [10] to circumvent the first limitation. This coder takes into account the geometry information at several key stages of the compression algorithm, i.e., intra prediction, block scanning and transform. This allows to drastically improve the compression performances. However, random access is not considered and the transmission of the whole image is still necessary. In this work, we propose a novel interactive coding scheme that is adapted to the random access of a user, and nevertheless takes into account the correlation between blocks by using some geometrical information. Instead of relying on a predictive coding approach [11], where the correlation exploitation requires the blocks processing/compression order to be fixed, we develop a coder based on the incremental entropy code [12] that can exploit these correlations, whatever the block order processing is. As a consequence, the block decoding order can be decided at the decoder side, once the request is known, such that only the visible blocks are transmitted. In other words, the proposed coder brings random access ability with a low storage cost compared to [10], but it significantly reduces the transmission rate. Even better, the transmission efficiency is as if the encoder already knew which part would be requested by the user, and can thus be considered as optimal.

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Fig. 1. Decomposition of a 3D model (a) into its texture atlas (b) and its triangular mesh (c).

II. ATLAS CODING WITH RANDOM ACCESS

In this section, we formalize the problem of atlas compression with random access at the decoder. As in most conventional image coders, the atlas is divided into blocks of dimension $B \times B$ that are represented as vectors $a_i$ of size
The atlas compression with random access is split into two steps. First, the atlas is compressed offline, i.e., it is coded once for all, independently from any user’s observation behavior. This offline encoding operation is denoted by a function \( f \). The entity \( f(A) \) is called the stored information, and the storage cost \( S \) is measured by \( S = |f(A)| \).

When observing a 3D model, only a part of the mesh, and thus of the texture, is visible. This is illustrated in Fig. 2. Let us denote by \( V \) the set of block indices that are visible by a user. Given the set \( V \), the second step of the compression is called the online extraction, and is denoted by \( g_V \). The entity \( g_V(f(A)) \) is called the extracted information and corresponds to an extractable subset of the stored information that is sufficient to recover the blocks whose index are in \( V \). The extraction cost is measured by \( E_V = |g_V(f(A))| \). Here, we emphasize that due to rendering delay constraint, decoding the atlas and reencoding the request can not be performed. The set of extractable information has to be carefully prepared in advance, and this is what makes the interactive compression challenging. For example, when the atlas is encoded with conventional image coders, there is no extractable information but the whole image itself. In other words, with conventional coders one has \( \forall V, E_V = S \). This looks suboptimal, and one can expect from an efficient interactive coder that \( E_V \leq S \). This is what we propose in the next section.

Once extracted, the decoder recovers a lossy version of the blocks \( \{\hat{a}_i\}_{i \in V} \). The coding distortion \( D_V \) is evaluated with a mean-squared error calculated in the rendered view domain between original view and the view generated with the lossy version of the atlas.

The goal of interactive coding is to minimize each of the elements of the triple \((S, E_V, D_V)\). Developing such a coder is particularly challenging for atlases. Indeed, minimizing \( S \) implies to jointly compress the \( a_i \), and this inevitably introduces dependencies between them, which makes their partial extraction nearly impossible. In the next section, we explain how we perform the joint compression of the \( a_i \) such that only a subset of them can be decoded.

III. PROPOSED INTERACTIVE CODER

In our proposed coder, as in the conventional ones, the blocks \( a_i \) are coded individually taking into account the correlation with the neighborhood. Let \( a_i \sim a_j \) denotes the fact that block \( a_i \) and \( a_j \) are adjacent. Two texture blocks are said to be adjacent if they correspond to two texture pieces that are side by side onto the surface of the 3D model. Let us define the neighborhood of the block with index \( i \) as the set of adjacent block indices, \( N(i) = \{ j \mid a_i \sim a_j \} \). In [10] as in all conventional coders, the set of blocks whose index are in \( N(i) \) and that are already coded are used to generate a prediction [11], denoted by \( \hat{a}_i \). A residue \( a_i - \hat{a}_i \) is then computed and coded with an arithmetic coder. This efficient compression strategy has however one main drawback: it imposes that the coding order at the encoder is the same as the coding order at the decoder, which contradicts random accessibility.

Instead, we propose to generate several predictions \( \hat{a}_i^k \) corresponding to different decoding orders that might appear at the decoder side, depending on user’s access. In fact, each \( \hat{a}_i^k \) corresponds to a prediction generated with blocks whose indices are in a given subset \( N^k(i) \subset N(i) \). Each subset \( N^k(i) \) corresponds to a possible decoding order, which depends on the user’s observation of the 3D model. Thanks to an incremental coder [12], we encode the block \( a_i \) based on the different predictions \( \hat{a}_i^k \) such that only the information, needed to complement the \( \hat{a}_i^k \) generated at the decoder, is extracted and sent to the decoder. Here we give more details on the different aforementioned steps. The proposed coder builds upon [10], and adds it the random accessibility at the decoder side. For the sake of clarity, some useful information from [10] are recalled in the following.

A. Neighborhood construction and prediction generation

Let us consider a block \( a_i \). The neighborhood construction consists in finding the set \( N(i) \) of the adjacent block’s indices. In a classical 2D image, \( N(i) \) is simply composed of the 4-connected neighbor indices placed on the previous/next
column/row. As discussed in our previous work [10], some textual information that are neighbors in the 3D model can be split into distant blocks in the atlas. Thanks to the mesh topology, we retrieve this proximity and describe it in $\mathcal{N}(i)$. This is illustrated in Fig. 3. At the decoder, the user decodes the set of blocks whose index are in $\mathcal{V}$. They are decoded in a given order $\sigma : \mathcal{V} \rightarrow [1,|\mathcal{V}|]$ (see Sec. III-C). For a given request $\mathcal{V}$, the set of blocks available for generating the predictions is $\mathcal{N}^\mathcal{V}(i) = \{ j \in \mathcal{N}(i) \mid \sigma(i) < \sigma(j) \}$, i.e., the set of blocks already decoded. Some reference pixels from the $\{a_j\}_{j \in \mathcal{N}^\mathcal{V}(i)}$ are put at the border of the block and forward propagated in an optimized directional order.

Considering all the requests $\mathcal{V}$, a set of $K$ scenarios are conceivable at the decoder (with $K << |\mathcal{V}|$ since many requests have a common $\mathcal{N}^\mathcal{V}(i)$). These scenarios lead to $K$ sets $\mathcal{N}^k(i)$ and thus $K$ predictions $\hat{a}_k^i$.

**B. Graph transform and incremental coding**

In order to compact the signal energy on a sparse basis, a transformation is then performed on each block $a_i$ and on each predicted block $\hat{a}_k^i$. Blocks of the texture image differ from classical 2D image block because a block may lie on a patch edge, and hence only a subset of pixels is informative. Therefore, it was proposed in [10] to perform a graph-based transform [13] in place of the regular 2D DCT.

One novelty of the proposed coder is to replace the conventional arithmetic coder by an incremental coding that can handle any prediction $\hat{a}_k^i$ at the decoder. More precisely, a codeword is built that can be seen as information (or parity bits) needed to correct the prediction and reconstruct the current block. Interestingly, this codeword is incremental. This means that the codeword can correct the worst prediction, but if the decoder has generated a better prediction, the sufficient bits to decode are extracted from the codeword. In [12] it was shown that theoretically the amount of bits sent to the decoder is the same as predictive coding, i.e. as if the navigation was known by the offline encoder. In practice, we use LDPC codes [14] applied on the quantized and binarized coefficients of the graph-based transform.

**C. Geometry-aware decoding order**

At the decoder side, only a subset $\mathcal{V}$ of the blocks are needed to be transmitted to the decoder. These blocks must be decoded in a smart order $\sigma$ such that when decoding a block $a_i$, the set of previously decoded neighboring blocks is as large as possible in order to generate the best possible prediction. Finding the optimal scanning order is however an NP-hard problem. We rather propose the approach illustrated in Fig. 4. The first block to be decoded is the access block (in red in Fig 4), which is one block coded in a standalone fashion, meaning that it is encoded independently and thereby can be decoded independently from its neighbors. The assumption, here, is that the navigation always starts at a position where this block is visible. This is without loss of generality because several access blocks can be introduced if several starting points are conceivable. Once this block is decoded, the whole patch where it belongs is decoded (in yellow in Fig 4), by exploring the blocks in a snake fashion so that there is always at least one neighboring block already decoded. Once this patch is entirely decoded, the algorithm explores the other patches where neighboring blocks are present (red arrows in Fig 4). Then neighboring patches are decoded and so forth, until all blocks are decoded.

For the later requests of the same user some blocks are already in memory. In that case, instead of starting by the access block, the decoding order begins with any of these blocks in memory and then proceeds as explained above.

**IV. EXPERIMENTAL VALIDATION**

In this experimental section, we compare our interactive coder with two baselines. The first one, called Whole Atlas (WA), encodes the whole atlas taking into account the geometry as in [10]. This method is very competitive for compression, and thus minimizes the storage $S$, but requires, for every user’s request, to transmit the entire atlas. Another baseline is called All Intra (AI), and corresponds to a method that is interactive, but is not able to take into account the correlation between the blocks. In other words, the method AI encodes each block independently as JPEG [2] does. For a fair comparison, the graph-based transform is performed in the method AI (as in ours and WA) by taking into account the geometry.

We used two 3D objects in our experiments, named Tombs [15] and Ewer [16], with atlas image size of 8192x8192 and 4096x4096 respectively. The objects are depicted in Fig. 5. We compress these two models with the three coding schemes leading to three storage costs: $S_{ours}$, $S_{WA}$ and $S_{AI}$. We also simulate three user navigations around each model, each last for 30 requests corresponding to successive changes of camera orientations and positions. At each instant we record the
In order to evaluate its impact, we rely on a comparison methodology proposed in [17], where the Bjontegaard metric is computed on the curve \((PSNR, R + \lambda S)\), for different values of \(\lambda\). Results are shown in Table I. We can see that when \(\lambda = 2\), meaning that storage matters more than transmission, the best method is the WA. This value corresponds to the size’s increase on the server. In other words, both All Intra and the proposed method have a small storage overhead, but for smaller \(\lambda\), \(i.e.,\) when a trade-off between storage and transmission is desired (transmission cost is becoming important), our method performs better than baselines.

V. CONCLUSION

In this paper, we proposed a new coder for mesh texture that enables interactivity. By replacing classical predictive coding scheme, by an incremental coding, and by adapting the intraprediction and the block scanning order, we have enabled our coder to transmit only what is requested by a user when looking at the 3D model from some points of view. The experiments demonstrate that only the useful information is sent to users and that the compression efficiency obtained by the non-interactive coders is maintained.

REFERENCES

TABLE I

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Fig. 6. transmission rate evolution over the user’s navigation for model Ewer at PSNR 38dB. (a) User 1. (b) User 2. (c) User 3.

Fig. 7. transmission rate evolution over the user’s navigation for model Tombs at PSNR 38dB. (a) User 1. (b) User 2. (c) User 3. Note that our proposed method is still better than all Intra method, but due to the scale of Y axis it is less visible.


