

Subprograms in “crème CAraMeL”

Lecture 9

Formal Languages and Compilers 2011

Nataliia Bielova

Assumptions and simplifications

- Let's add
 - declarations of subprograms (procedure)
 - execution of subprograms (call and passing the parameters)
- no declarations inside the blocks `begin ... end`
- declarations are non-static: dynamic local environment
- only one environment, no activation record: dynamic non local environment
- passing the parameters only *by value*

Procedure: example

```
program
  var x : int

  procedure proc1 (a: int)
  begin
    write(a)
  end;

  procedure proc()
  var x : int
  begin
    x:=5;
    call proc1(x);
    write(x)
  end

begin
  x := 40;
  write(x);
  call proc();
  call proc1(x)
end
```



```
40
5
5
40
```

Procedure: implementation

Syntax:

- parser.mly: new token PROCEDURE, CALL, COMMA (“,”)
- lexer.mll: strings corresponding to token
- syntaxtree.ml: constructors for
 - declarations
 - formal parameters
 - calls
 - actual parameters
 - other modifications
- parser.mly: productions to construct new nodes

Procedure: implementation

Semantics:

- new value in the environment:

`Descr_Procedure of param list * dec list * cmd`

- declaration
- execution (call)
 - evaluation of actual parameters
 - type checking for the list of parameters
 - actual execution of the procedure

Function: example

```
program
  var x : int

  function fact(a: int): int
    var b : int
    begin
      if (a = 0) then
        fact := 1
      else begin
        b := call fact(a - 1);
        fact := a * b
      end
    end
  end

begin
  x := call fact(12);
  write(x)
end
```



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Function: implementation

- Syntax:
 - keyword: function
 - new nodes for: declaration, execution (call), evaluation (call!)
 - adjust the syntax tree
- Semantics:
 - declaration (attention: a location for return value is needed!)
 - evaluation
 - execution

Projects for the exam

Extension of an interpreter/compiler:

- 1 person (either one or another item)
 - pointers and dynamic memory
 - pointers and record
 - vectors “by linked list”
- 2 persons (also here)
 - multidimensional matrices and horizontal-vertical slices
 - pointers, passing parameters by value, name, reference, value-result
- 3 persons (like before)
 - record, pointers, multidimensional matrices and horizontal-vertical slices
 - definition of functions and nested procedures (static scoping) and passing parameters by value, name, reference, value-reference

Pointers

- Declaration:

```
var p : ^int;
```

```
var q: ^^float;
```

- Referencing(@) and dereferencing (^) :

```
x := 1;
```

```
p := @x;
```

```
y := ^p + 4;
```

If x and y are integers, then in the end $y = 5$.

Dynamic memory

Add to the language the possibility to use pointers and allocation/deallocation of dynamic memory (heap), using *one* of the following approaches of memory release:

- reference counter
- garbage collection

A correct implementation will allow to create and use the dynamic data structures using pointers in crème CAraMeL.

Vectors “by linked list”

Change the implementation of vectors in crème CAraMeL in a way that the following operations are possible:

$v(i) := 5$ inserting an element (growing the length of the vector)

$v[i] := 5$ substitution of an element (the length remains the same)

$v\#i$ deleting an element (the vector becomes shorter)

$v?5$ returns an integer i if vector contains value 5 at position i and an integer -1 if there is no value 5 in the vector

Record

- Definition:

```
type name_record = record {  
    name_field1 : type;  
    ...  
    name_fieldn : type;  
}
```

- Declaration:

```
var v : name_record;
```

- Access:

```
v.name_fieldi := expression;  
a := v.name_fieldi;
```