# Customizing Fixed-Point and Floating-Point Arithmetic – A Case Study in K-Means Clustering

Benjamin Barrois and Olivier Sentieys

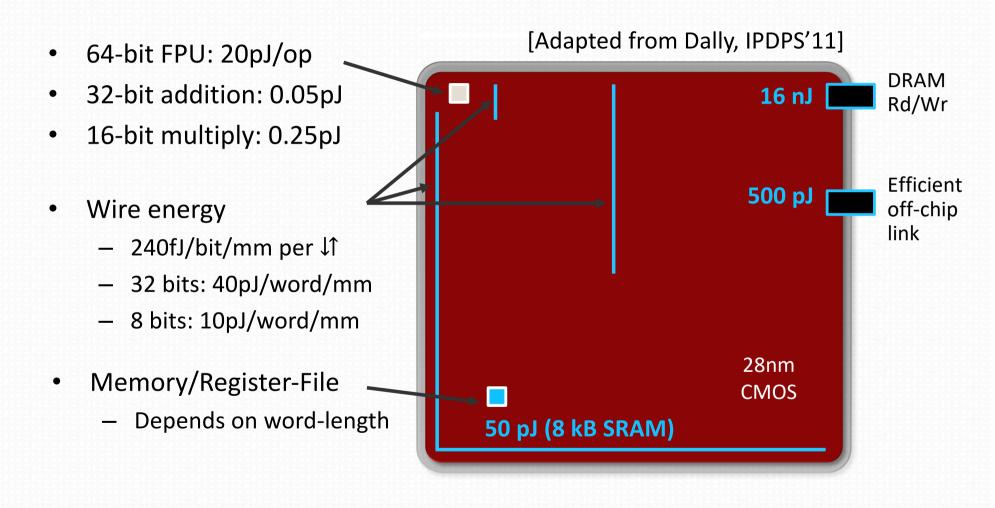
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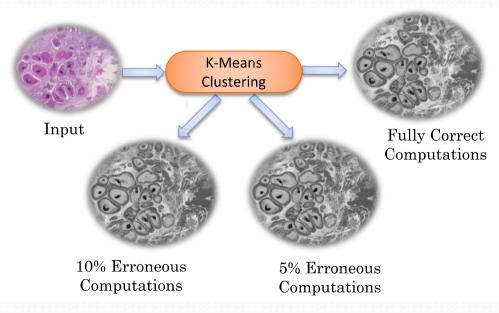
### Energy Cost in a Processor/SoC



Energy strongly depends on data representation and size

### Many Applications are Error Resilient

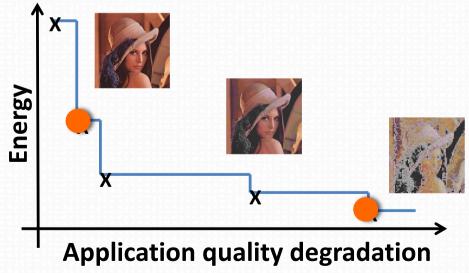
- Produce outputs of acceptable quality despite approximate computation
  - Perceptual limitations
  - Redundancy in data and/or computations
  - Noisy inputs
- Digital communications, media processing, data mining, machine learning, web search, ...



## **Approximate Computing**

- Play with number representation to reduce energy and increase execution speed while keeping accuracy in acceptable limits
  - Relaxing the need for fully precise operations

- Trade quality against performance/energy
  - Compile-time/run-time
- Different levels
  - Operators/functions/algorithms



#### Outline

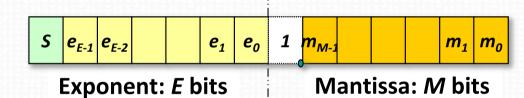
- Introduction
- Number Representation
  - Fixed-Point
  - Floating-Point
- Customizing Arithmetic Operators
- Direct Comparison of Custom Operators
- ApxPerf Framework
- Results on K-Means Clustering Algorithm
- Conclusions

### Number Representation

Floating-Point (FIP)

$$x = (-1)^s \times m \times 2^{e-127}$$

s: sign, m: mantissa, e: exponent



- Easy to use
- High dynamic range
- IEEE 754

| Format                  | е  | m  | bias |
|-------------------------|----|----|------|
| Single Precision        | 8  | 23 | 127  |
| <b>Double Precision</b> | 11 | 52 | 1023 |

Fixed-Point (FxP)

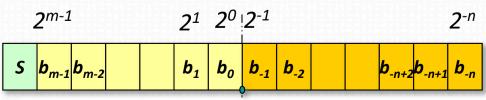
$$x = p \times K$$

p: integer,  $K=2^{-n}$ : fixed scale factor

- Integer arithmetic
- Efficient operators
  - Speed, power, cost

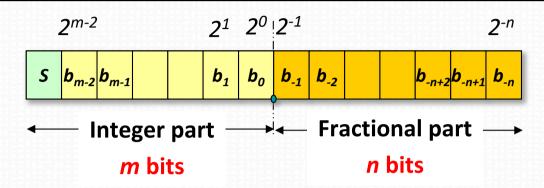
$$x = s.(-2)^m + \sum_{i=-n}^{m-1} b_i.2^i$$

s: sign, m: magnitude, n: fractional



Integer part: m bits Fractional part: n bits

#### Fixed-Point Arithmetic



$$x = s.(-2)^{m-1} + \sum_{i=-n}^{m-2} b_i.2^i$$

s: sign, m: magnitude, n: fractional

Accuracy (error)

Dynamic range

1.640625  
+ 2.5125  
= 4.153125 
$$x \in [-4; 4[$$
  
overflow

Need for explicit normalization

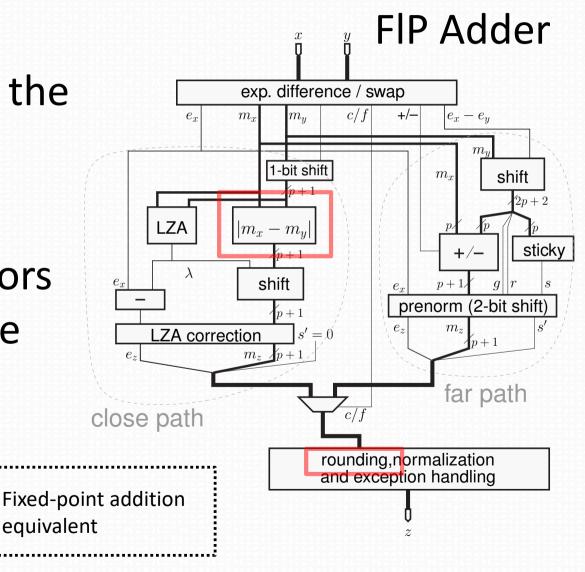
$$(int)(((INT64)a * (INT64)b) >> N)$$

Use of popular libraries (e.g. sc\_fixed, ac\_fixed)

### Floating-Point Arithmetic

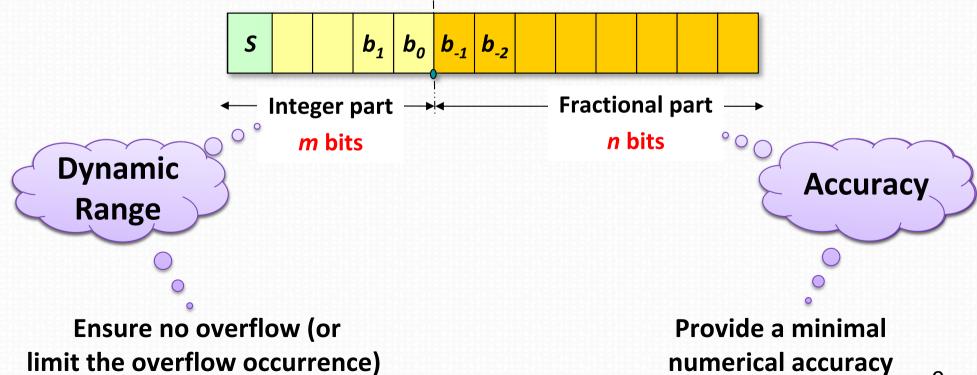
 Floating-point hardware is doing the job for you!

 Arithmetic operators are therefore more complex



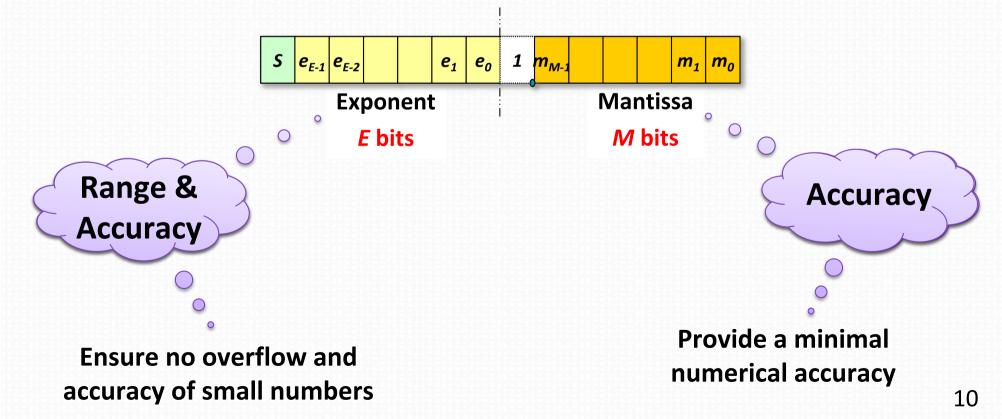
## **Customizing Fixed-Point**

- Minimize word-length W=m+n
- Determine integer and fractional parts
  - Fixed-point refinement



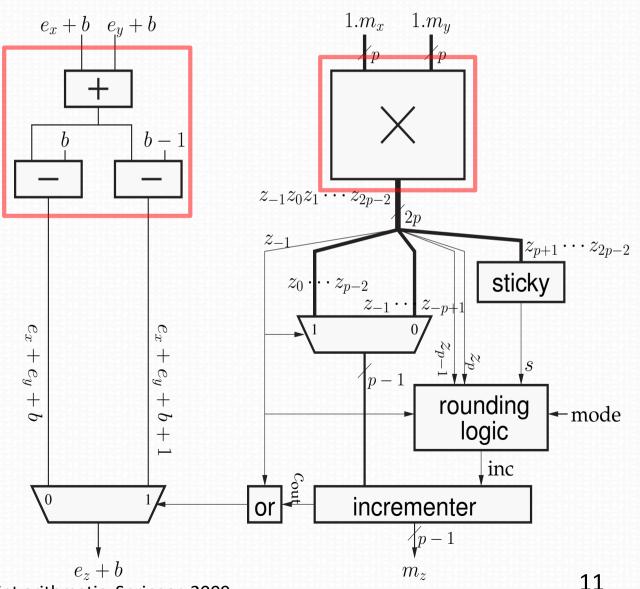
### **Customizing Floating-Point**

- Minimize word-length W=E+M+1
- Determine exponent and mantissa (and bias)
- Error is relative to number value



### Floating-Point Multiplication

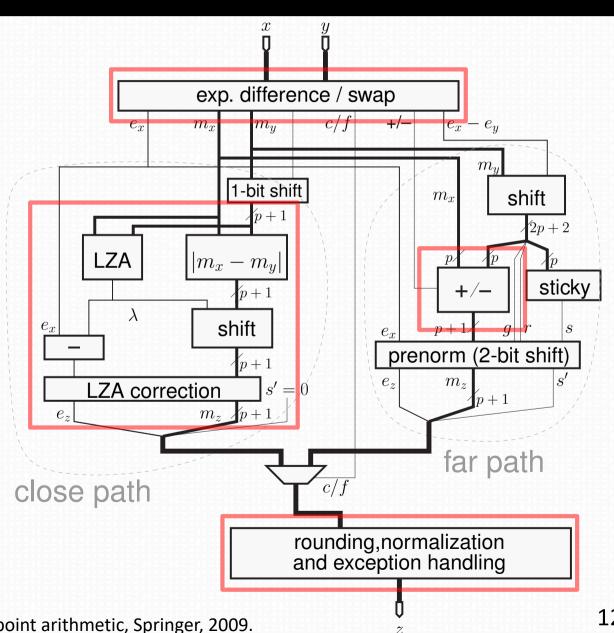
- Representation (W,E,M)
- Exponent e
  - E bits
- Mantissa m
  - M bits



J.-M. Muller et al., Handbook of Floating-point arithmetic, Springer, 2009.

### Floating-Point Addition

- Representation (W,E,M)
- Exponent e
  - E bits
- Mantissa m
  - M bits



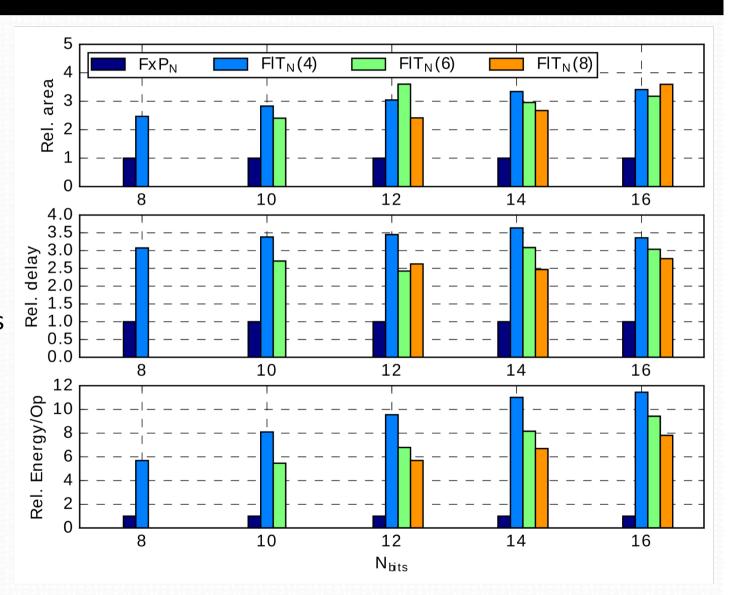
### ct\_float: a Custom-FIP C++ Library

- ct\_float: a Custom Floating-Point C++ Library
  - Operator simulation and (High-Level) synthesis
  - Templated C++ class
    - Exponent width e (int)
    - Mantissa width m (int)
    - Rounding method r (CT\_RD,CT\_RU,CT\_RND,CT\_RNU)
  - Many synthetizable overloaded operators
    - Comparison, arithmetic, shifting, etc.

```
ct_float<8,12,CT_RD> x,y,z;
x = 1.5565e-2;
z = x + y;
```

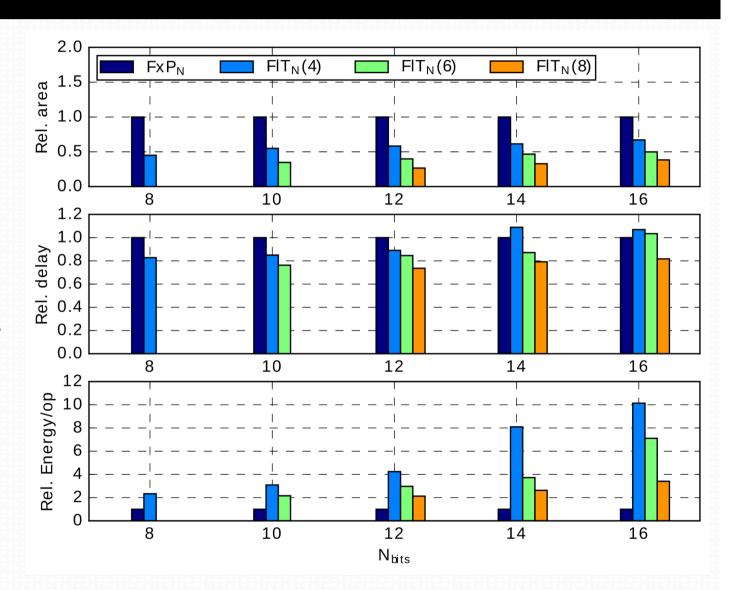
### FxP vs. FIP: Adders

- FxP<sub>N</sub>
  - Fixed-Point
  - N bits
- FIT<sub>N</sub>(E)
  - Floating-Point
  - N bits
  - Exponent *E* bits
- FxP adders are always smaller, faster, less energy



### FxP vs. FIP: Multipliers

- FxP<sub>N</sub>
  - Fixed-Point
  - N bits
- FIT<sub>N</sub>(E)
  - Floating-Point
  - N bits
  - Exponent E bits
- FIP multipliers
   are smaller,
   faster, but more
   consuming

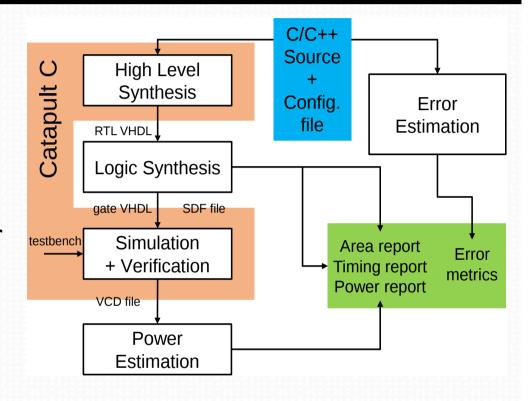


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### **Energy-Accuracy Trade-offs**

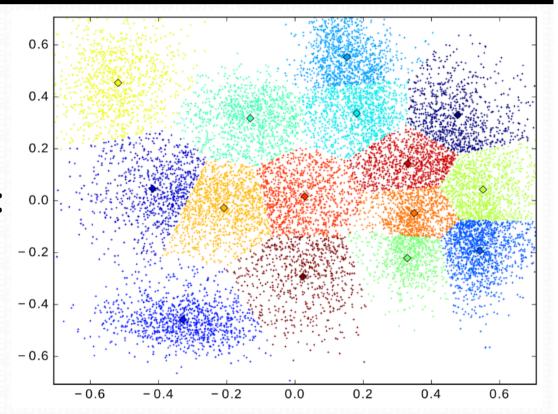
- ApxPerf2.0 framework
  - Based on C++ templates,
     HLS, and Python
  - VHDL and C/C++ operator descriptions
    - Approximate, FxP, FIP



- Fully automated characterization
- Generates delay, area, and power results
- Extract error metrics
  - mean square error, mean average error, relative error, min/max error, bit error rate, etc.

### K-Means Clustering

- Data mining, image classification, etc.
- A multidimensional space is organized as:
  - -k clusters  $S_i$ ,
  - $-S_i$  defined by its centroid  $\mu_i$



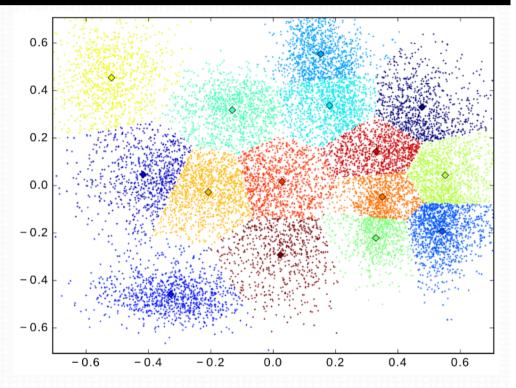
• Finding the set of clusters  $S = \{S_i\}_{i \in [0,k-1]}$ 

$$\underset{S}{\operatorname{arg\,min}} \sum_{i=1}^{k} \sum_{s \in S} \|x - \mu_i\|^2$$

satisfying

### K-Means Clustering

- Lloyd's iterative algorithm
  - approximations of the optimal centroids
  - estimation-maximization three-step iterative process



- Distance computation  $d \leftarrow (x y) \times (x y)$
- Iteration of computations until
  - sum of distances from data points x to centroid  $\mu_i$  between two iterations is less than a given threshold
  - maximum number of iterations

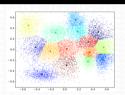
#### Experimental setup

- 20 data sets composed of 15.10<sup>3</sup> samples
- Gaussian distributions with random covariance matrices around 15 random mean points
- Accuracy targets:  $10^{-2}$ ,  $10^{-3}$ ,  $10^{-4}$
- Reference is double-precision floating-point
- 28nm, 100MHz
- Error metrics
  - Mean square error of cluster centroids (CMSE)
    - lower is better
  - Classification error rate (ER)
    - i.e. proportion of points not being tagged by cluster identifier

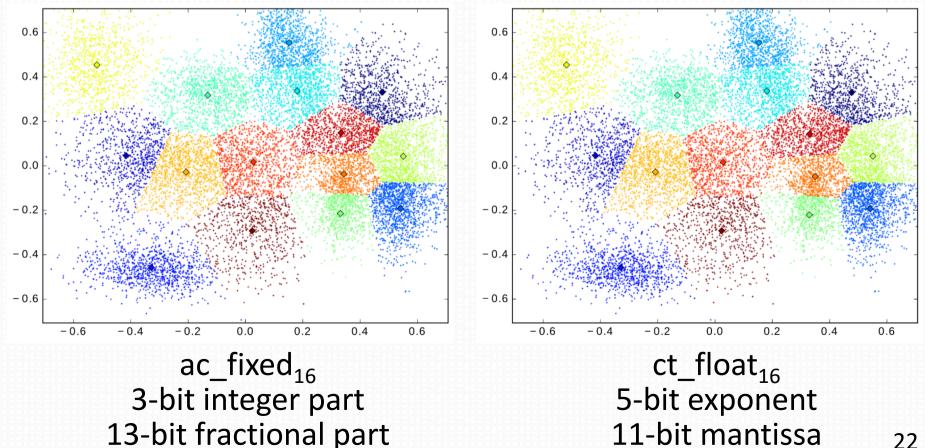
- Results with 8-bit and 16-bit FIP and FxP arithmetic operators
- Stopping condition set to 10<sup>-4</sup>

|                          | $\operatorname{ct\_float}_8(5)$ | $\operatorname{ct\_float}_{16}(5)$ | $ac_fixed_8(3)$ | $ac_{16}(3)$ |
|--------------------------|---------------------------------|------------------------------------|-----------------|--------------|
| Area $(\mu m^2)$         | 392.3                           | 1148                               | 180.7           | 575.1        |
| $N_{ m cycles}$          | 3                               | 3                                  | 2               | 2            |
| $E_{\rm dc} (nJ)$        | 1.23E-4                         | 5.99E-4                            | 5.03E - 5       | 3.25E-4      |
| $N_{ m it}$              | 8.35                            | 59.3                               | 14.9            | 65.1         |
| $E_{\text{K-means}}(nJ)$ | 38.24                           | 1100                               | 23.90           | 644.34       |
| CMSE                     | 1.75E - 3                       | 3.03E-7                            | 1.85E-2         | 3.28E-7      |
| Error Rate               | 35.1 %                          | 2.94 %                             | 62.3 %          | 0.643 %      |

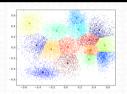
• W = 16 bits, accuracy =  $10^{-4}$ 



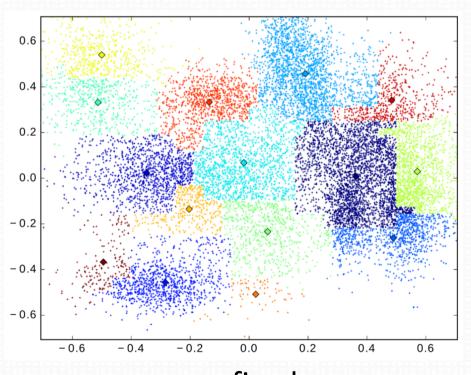
No major (visible) difference with reference



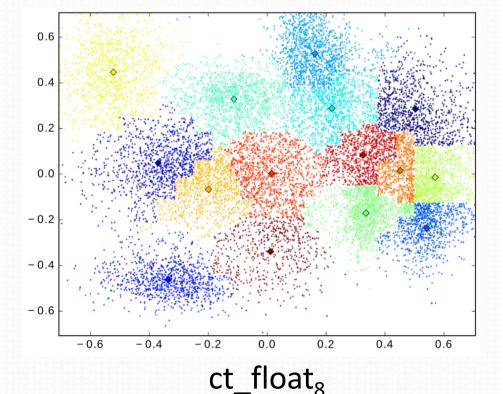
• W = 8 bits, accuracy =  $10^{-4}$ 



8-bit float is better and still practical



ac\_fixed<sub>8</sub>
3-bit integer part
5-bit fractional part

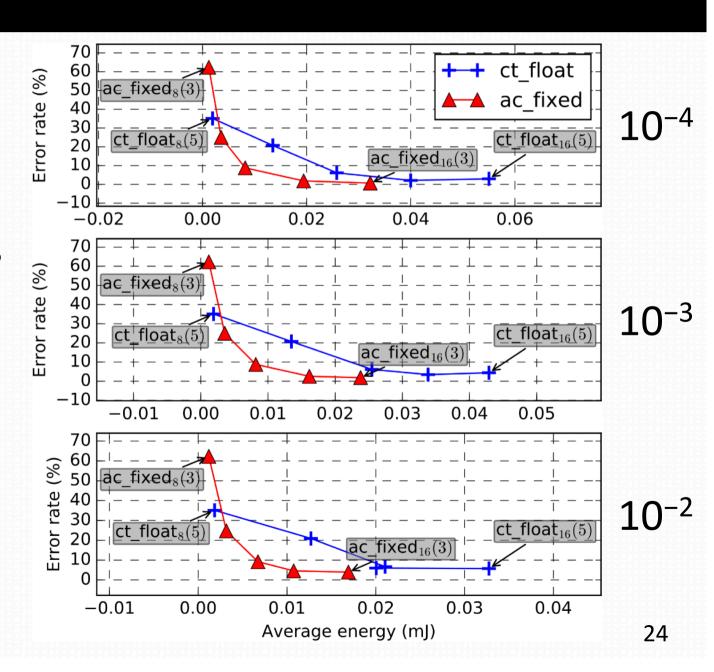


5-bit exponent 3-bit mantissa

### **Energy versus Classification Error Rate**

Average energy consumed by K-means algorithm

 Stopping conditions:
 10<sup>-2</sup> to 10<sup>-4</sup>

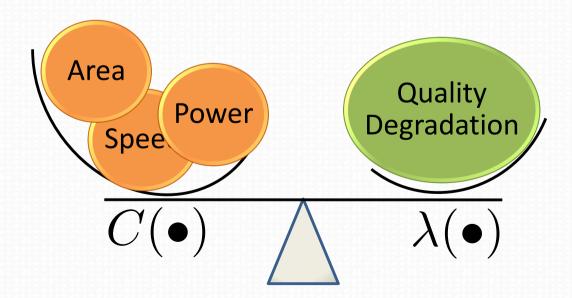


#### Conclusions

- Total energy (algorithm) depends on:
  - Energy of arithmetic operations
  - Algorithm convergence speed
- Slower increase of errors for floating-point
- Small floating-point (e.g. 8-bit) provides better error rate/energy ratio
- Perspectives
  - Custom exponent bias in ct\_float
  - Towards an automatic optimizing compiler considering both FxP and FIP representations

### Customizing Number Representation

- Loss of accuracy incurs quality degradation
- Essentially, an optimization process
  - Determine the number of bits for each data
  - Determine the format for each data



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