

# Elliptic curves for SNARK and proof systems

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25th Workshop on Elliptic Curve Cryptography  
Taipei, Taiwan, October 30, 2024

joint work with Simon Masson

These slides at <https://people.rennes.inria.fr/Aurore.Guillevic/talks/2024-10-ECC/24-10-30-ECC-Aurore.pdf>

zk-SNARK

Elliptic Curves and Pairings

Proof-friendly curves

More embedded curves

More families of embedded curves

# Outline

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# Zero-knowledge proofs (ZKP)

**Alice**

I know the solution to this complex equation

Examples:

On this chess board, I know mat in 3 moves

I know where is Wally/Waldo on this drawing

I know a solution to this sudoku grid

I know a preimage of this hash function value

**Bob**

No idea what the solution is but Alice claims to know it



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- **Sound:** **Alice** has a **wrong solution**  $\implies$  **Bob** is **not convinced**.

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- **Complete:** Alice has the solution  $\implies$  Bob is convinced.

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- **Sound:** Alice has a wrong solution  $\implies$  Bob is not convinced.
- **Complete:** Alice has the solution  $\implies$  Bob is convinced.
- **Zero-knowledge:** Bob does NOT learn the solution.

## Example: Sigma protocol

Alice

I know  $x \in \mathbf{Z}_q$  such that  
 $g^x = y$  in  $\mathbf{G}$ ,  $\#\mathbf{G} = q$  prime

Bob



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$$n \xleftarrow{\$} \mathbf{Z}_q$$

$$A = g^n$$

Bob

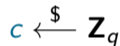
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$$n \xleftarrow{\$} \mathbf{Z}_q \quad \xrightarrow{A = g^n}$$

$$\xleftarrow{c}$$

$$s = n + c \cdot x \pmod q \quad \xrightarrow{s}$$

Bob

$$c \xleftarrow{\$} \mathbf{Z}_q$$

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Bob

$$c \xleftarrow{\$} \mathbf{Z}_q$$

$$g^s \stackrel{?}{=} A \cdot y^c$$

$$\text{with } A \cdot y^c = g^n \cdot g^{x \cdot c}$$

$$\text{then } g^n \cdot g^{x \cdot c} = g^{n+x \cdot c}$$

Hide the verification  
in the exponents  
(the scalar field)

# Non-Interactive Zero-Knowledge (NIZK) Sigma protocol

Alice

Bob

I know  $x$  such that  $g^x = y$

$\mathbf{G}, g, y$

$$\begin{aligned} n &\stackrel{\$}{\leftarrow} \mathbf{Z}_q, A = g^n \\ c &= H(A, y) \\ s &= n + c \cdot x \pmod q \end{aligned}$$

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# Non-Interactive Zero-Knowledge (NIZK) Sigma protocol

Alice

I know  $x$  such that  $g^x = y$

$\mathbf{G}, g, y$

Setup

$$n \xleftarrow{\$} \mathbf{Z}_q, A = g^n$$

$$c = H(A, y)$$

$$s = n + c \cdot x \pmod q$$

Prove



$$\pi = (A, c, s)$$

proof

Bob

$$g^s \stackrel{?}{=} A \cdot y^c$$

$$c \stackrel{?}{=} H(A, y)$$

Verify



# zk-SNARK: Zero-Knowledge Succinct Non-interactive ARgument of Knowledge

"I have a *computationally sound, complete, zero-knowledge, succinct, non-interactive* proof that a statement is true and that I know a related secret".

## Succinct

A proof is very *short* and *easy* to verify.

## Non-interactive

No interaction between the prover and verifier for proof generation and verification (except the proof message).

## ARgument of Knowledge

Honest verifier is convinced that a computationally bounded prover knows a secret information.

# zk-SNARKs in a nutshell

## Main ideas:

1. Reduce a **general statement** satisfiability to a polynomial equation satisfiability.
2. Use Schwartz–Zippel lemma to succinctly verify the polynomial equation with high probability.
3. Use homomorphic hiding cryptography to blindly verify the polynomial equation.
4. Make the protocol non-interactive.

# Needs of groups for proof systems and SNARK

## Statement

group  $\mathbf{G}'$  of prime order over  $\mathbb{F}_q$  /  
Hash function over base field  $\mathbb{F}_q$

- ed\_25519 signature verification  
 $q = 2^{255} - 19$
- Hash function verification  $y = H(x)$   
 $H$ : Poseidon, Anemoi...

## Proof

group  $\mathbf{G}$  of prime order  $q$  over  $\mathbb{F}_p$

Group where multiplication  
in the exponents is possible:  
given  $g^a, g^b$ , compute  $g^{ab}$   
without knowing  $a, b$   
 $\rightarrow \approx$  pairing-friendly curves

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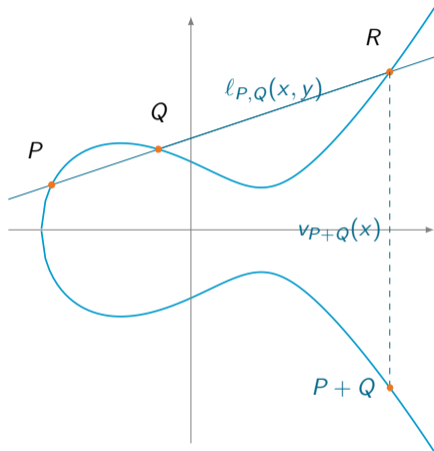
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# Elliptic curve $E/\mathbb{F}_p: y^2 = x^3 + ax + b$ , $a, b \in \mathbb{F}_p$ , $p \geq 5$ , group law



- $E(\mathbb{F}_p)$  has an efficient group law  $\rightarrow \mathbf{G}_1$  (chord and tangent rule)
- $\#E(\mathbb{F}_p) = p + 1 - t$ , trace  $t: |t| \leq 2\sqrt{p}$
- large prime  $q \mid p + 1 - t$  coprime to  $p$
- $E(\mathbb{F}_p)[q] = \{P \in E(\mathbb{F}_p) : [q]P = \mathcal{O}\}$  has order  $q$
- $E[q] \simeq \mathbf{Z}/q\mathbf{Z} \times \mathbf{Z}/q\mathbf{Z}$  (for crypto)
- only generic attacks against DLP on well-chosen genus 1 and genus 2 curves
- optimal parameter sizes

## Pairing as a black box

$(\mathbf{G}_1, +)$ ,  $(\mathbf{G}_2, +)$ ,  $(\mathbf{G}_T, \cdot)$  three cyclic groups of large prime order  $q$

Pairing: map  $e : \mathbf{G}_1 \times \mathbf{G}_2 \rightarrow \mathbf{G}_T$

1. bilinear:  $e(P_1 + P_2, Q) = e(P_1, Q) \cdot e(P_2, Q)$ ,  $e(P, Q_1 + Q_2) = e(P, Q_1) \cdot e(P, Q_2)$
2. non-degenerate:  $e(G_1, G_2) \neq 1$  for  $\langle G_1 \rangle = \mathbf{G}_1$ ,  $\langle G_2 \rangle = \mathbf{G}_2$
3. efficiently computable.

Most often used in practice: swap scalars, multiply in the exponents

$$e([a]P, [b]Q) = e([b]P, [a]Q) = e(P, Q)^{ab} .$$

Can multiply only once!

$\rightsquigarrow$  Many applications in asymmetric cryptography.

# Cryptographic pairing

## Modified Weil or Tate pairing on an elliptic curve


Discrete logarithm problem with one more dimension.

$$e : E(\mathbb{F}_p)[q] \times E(\mathbb{F}_{p^k})[q] \longrightarrow \mathbb{F}_{p^k}^*, \quad e([a]P, [b]Q) = e(P, Q)^{ab}$$

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### Attacks

- inversion of  $e$  : hard problem (exponential)
- discrete logarithm computation in  $E(\mathbb{F}_p)$  : hard problem (exponential, in  $O(\sqrt{q})$ )
- discrete logarithm computation in  $\mathbb{F}_{p^k}^*$  : **easier, subexponential**  $\rightarrow$  take a large enough field



## Finding pairing-friendly curves

Designed on purpose: otherwise  $k \approx q$

Choose prime integer  $q$ , degree  $k$  then obtain  $p$ : inefficient curve

Design families: parameterized  $p(x)$ ,  $q(x)$ ,  $t(x)$

- Complex Multiplication (CM) equation:  $t^2 - 4p = -Dy^2$
- (Compute  $t^2 - 4p$ , get its square-free factorization)
- $D$  discriminant, square-free (in number theory, if  $D = 1, 2 \pmod{4}$  then  $D \leftarrow 4D$ )

SEA: from coefficients to parameters

$$E/\mathbb{F}_p: y^2 = x^3 + ax + b$$

Schroof–Elkies–Atkin (SEA)

compute trace  $t$

order  $q = p + 1 - t$

iterate over  $a, b$  until  $q$  is prime

CM: from parameters to coefficients

base field  $\mathbb{F}_p$ , trace  $t$ , order  $q$

CM equation  $t^2 - 4p = -Dy^2$

compute Hilbert Class polynomial  $H_D(X)$

compute a root  $j$ ,  $H_D(j) = 0 \pmod{p}$

$$j(E) = 1728 \frac{4a^3}{4a^3 + 27b^2}$$

$$E/\mathbb{F}_p: y^2 = x^3 + \frac{3j}{j-1728}x + \frac{2j}{1728-j}$$

## First ordinary pairing-friendly curves: MNT [MNT01]

Miyaji, Nakabayashi, Takano,  $\#E(\mathbb{F}_p) = p(x) + 1 - t(x) = q(x)$

$$k = 3 \begin{cases} t(x) = -1 \pm 6x \\ q(x) = 12x^2 \mp 6x + 1 \\ p(x) = 12x^2 - 1 \\ Dy^2 = 12x^2 \pm 12x - 5 \end{cases}$$

$$k = 4 \begin{cases} t(x) = -x, x + 1 \\ q(x) = x^2 + 2x + 2, x^2 + 1 \\ p(x) = x^2 + x + 1 \\ Dy^2 = 3x^2 + 4x + 4 \end{cases}$$

$$k = 6 \begin{cases} t(x) = 1 \pm 2x \\ q(x) = 4x^2 \mp 2x + 1 \\ p(x) = 4x^2 + 1 \\ Dy^2 = 12x^2 - 4x + 3 \end{cases}$$

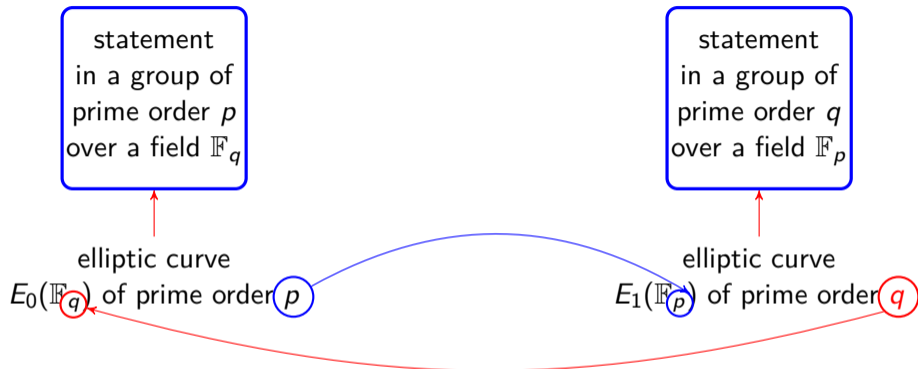
CODA [MS18]:

$k = 6$ , 753 bits,  $E_6 \approx 137$  bits of security,  $D = -241873351932854907$ , seed  $u =$

0xaa3a58eb20d1fec36e5e772ee6d3ff28c296465f137300399db8a5521e18d33581a262716214583d3b89820dd0c000

$k = 4$ , 753 bits,  $E_4 \approx 113$  bits of security

## Cycle of curves: unlimited chains of SNARKs [BCTV14]



MNT-4 and MNT-6 curves form a cycle

$k = 4$ , MNT-4:  $t = -x$ , order  $q = x^2 + 1$ , field  $p = x^2 + x + 1$

$k = 6$ , MNT-6 ( $x \leftrightarrow 2x$ ):  $t' = 1 - x$ , order  $p = x^2 + x + 1$ , field  $q = x^2 + 1$

Unique known cycle of pairing-friendly curves. Impossibility results: [CCW19, BMUS23]

New constructions with higher dimensional curves [CCN24, CK24]

## Very popular pairing-friendly curves: Barreto-Naehrig (BN) [BN06]

$$E_{BN} : y^2 = x^3 + b, \quad p \equiv 1 \pmod{3}, \quad D = 3 \text{ (ordinary)}, \quad j_E = 0$$

$$p = 36u^4 + 36u^3 + 24u^2 + 6u + 1$$

$$t = 6u^2 + 1$$

$$q = p + 1 - t = 36u^4 + 36u^3 + 18u^2 + 6u + 1$$

$$t^2 - 4p = -3(6u^2 + 4u + 1)^2 \rightarrow \text{no CM method needed}$$

Comes from the Aurifeuillean factorization of  $\Phi_{12} : \Phi_{12}(6u^2) = q(u)q(-u)$

Security level	$\log_2 q$	$\log_2 p$	$k$	finite field	$\rho = \log p / \log q$
102	256	256	12	3072	1
123	384	384	12	4608	1
132	448	448	12	5376	1

Formerly BN-254 in Euthereum with seed 0x44e992b44a6909f1

## Barreto, Lynn, Scott curves [BLS03]

Any  $k$ ,  $3 \mid k$ ,  $18 \nmid k$  possible

BLS12 ( $k = 12$ ) becomes more and more popular, replacing BN curves

$$E_{\text{BLS}} : y^2 = x^3 + b, \quad p \equiv 1 \pmod{3}, \quad D = 3 \text{ (ordinary)}$$

$$p = (u - 1)^2 / 3(u^4 - u^2 + 1) + u$$

$$t = u + 1$$

$$q = (u^4 - u^2 + 1) = \Phi_{12}(u)$$

$$p + 1 - t = \underbrace{(u - 1)^2 / 3(u^4 - u^2 + 1)}_{\text{cofactor}}$$

$$t^2 - 4p = -3y(u)^2 \rightarrow \text{no CM method needed}$$

BLS12-381 (Zcash [Bow17]) with seed `-0xd201000000010000`

BLS12-377 (Zexe [BCG<sup>+</sup>]) with seed `0x8508c00000000001`

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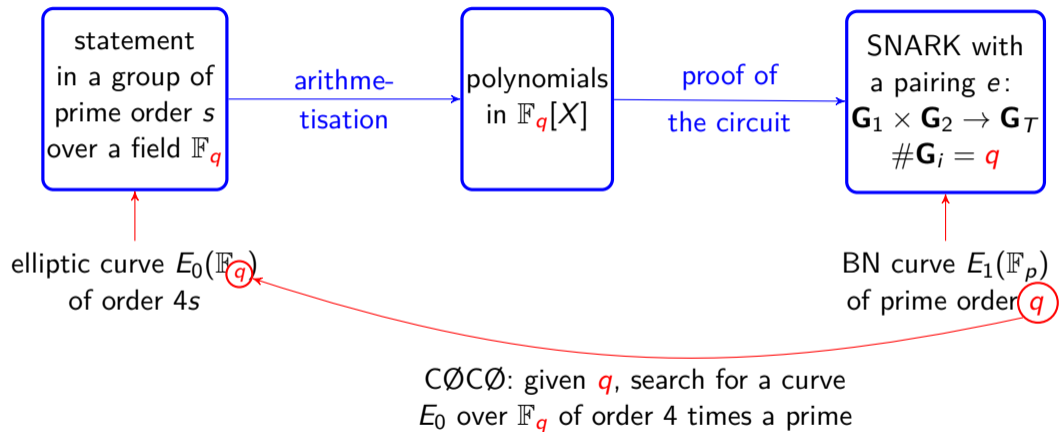
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# $C\emptyset C\emptyset$ embedded curve: Kosba et al. construction [KZM<sup>+</sup>15]



## Embedded SNARK-friendly curves

Usually a twist-secure elliptic curve in Montgomery or (twisted) Edwards form

Input: base field  $\mathbb{F}_q$

Output: an embedded curve over  $\mathbb{F}_q$  of order  $4s$  or  $8s$  with prime  $s$

Procedure: Increment the curve coefficients until a suitable curve is found  
(Nothing-up-my-sleeves strategy)

CØCØ [KZM<sup>+</sup>15] with BN-254a,

JubJub [ZCa21] or Bandersnatch [MSZ21] with BLS12-381,

first attempt to generalize Bandersnatch [SEH24]



## Bandersnatch [MSZ21]

- Find an embedded elliptic curve  $E'$  over  $\mathbb{F}_{q_{\text{BLS12-381}}}$  of trace  $t'$ , above BLS12-381
- With a small discriminant  $D'$  in  $t'^2 - 4q = -D'y'^2$  to allow faster scalar multiplication with GLV
- twist-secure:  $q + 1 - t'$ ,  $q + 1 + t'$  contain a large prime
- Use the CM method

$u = -0\text{xd}201000000010000$ ,  $q = u^4 - u^2 + 1$  is prime (BLS12-381)

The trace  $t'$  can be any integer in the range  $(-2\sqrt{q}; 2\sqrt{q})$

Idea: enumerate small  $D'$ , get  $t'$ , order  $s$ , twist order  $s'$  until  $s, s'$  contain a large prime

Bandersnatch curve:  $D' = 2$  (i.e.  $D' = -8$ ),  $s = 2^2 \times p_{253}$ ,  $s' = 2^7 \cdot 3^3 \times p_{244}$

Is it a *magical* curve? It is *too good to be true*?

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## More Bandersnatch curves

Extend the search space for discriminants  $D'$

Rewrite the algorithm to enumerate the curves much faster

We get more embedded twist-secure curves with BLS12-381:

- $D = 2$ , Bandersnatch
- $D = 1030258$ ,  $r = 4p_{253}$ ,  $r' = 2^3 \cdot 7p_{250}$
- $D = 1429201$ ,  $r = 4p_{253}$ ,  $r' = 2^8 \cdot 5p_{245}$
- $D = 1470074$ ,  $r = 2^9 p_{246}$ ,  $r' = 2^2 \cdot 3^4 \cdot 5p_{245}$
- $D = 1992138$ ,  $r = 2^7 p_{248}$ ,  $r' = 2^2 \cdot 3^2 \cdot 79p_{244}$
- $D = 7636102$ ,  $r = 2^2 p_{253}$ ,  $r' = 2^3 \cdot 3^2 \cdot 23p_{245}$
- ...

More embedded prime-order curves with BLS12-381:

- $D = 6673027$ ,  $r$  prime,  $r' = c \cdot p_{234}$  (twist-secure)
- $D = 7321939$ ,  $r$  prime,  $r' = c' \cdot p_{206}$

# Imaginary Quadratic Number Field

Let  $d > 0$  a square-free integer.

$$K = \mathbf{Q}[x]/(x^2 + d) \simeq \mathbf{Q}(\sqrt{-d})$$

is a imaginary quadratic number field whose maximal order is  $\mathcal{O}_K = \mathbf{Z}[\tau]$  where

$$\tau = \begin{cases} \sqrt{-d} & \text{if } d \not\equiv 3 \pmod{4}, \\ \frac{1+\sqrt{-d}}{2} & \text{if } d \equiv 3 \pmod{4}. \end{cases}$$

The **norm** of an algebraic integer  $a + b\tau$ ,  $a, b \in \mathbf{Z}$  is

$$N_{K/\mathbf{Q}}(a + b\tau) = \begin{cases} \text{Res}(a + bX, X^2 + d) = a^2 + db^2 & \text{if } d \equiv 1, 2 \pmod{4} \\ \text{Res}(a + bX, X^2 - X + \frac{d+1}{4}) = a^2 + ab + \frac{d+1}{4}b^2 & \text{if } d \equiv 3 \pmod{4}. \end{cases}$$

## Solving norm equation

Given positive integer  $n$ , find  $\eta \in \mathbf{Z}[\tau]$  of norm  $n \rightarrow$  sometimes no solution

Example:  $K = \mathbf{Q}(\sqrt{-5})$ ,  $\tau = \sqrt{-5}$ . Solve for  $\pi = a + b\tau \in \mathcal{O}_K$ ,

$$N_{K/\mathbf{Q}}(\pi) = a^2 + 5b^2 = p \text{ prime}$$

- ramified primes  $p = 2, 5$
- inert primes 11, 13, 17, 19, 31, 37, 53, 59, 71, 73, 79, 97
- splitting primes 3, 7, 23, 43, 47, 67, 83 but no solution
- splitting primes having solutions:  $29 = N(3 \pm 2\tau)$ ,  $41 = N(6 \pm \tau)$ ,  
 $61 = N(4 \pm 3\tau)$ ,  $89 = N(3 \pm 4\tau)$

In average (over  $10^5$   $p$ ):  $1/2$  are split,  $1/4$  have a solution  $\pi$ ,  $p = N(\pi)$ .

- $-d$  is a square in half the cases
- prime ideal  $\mathfrak{p}$  above  $p$  is principal with  $1/h(K)$  chance,  $h(K) =$  **Class Number**  
then  $\pi = a + b\tau$  exists,  $N_{K/\mathbf{Q}}(\pi) = p$

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**Algorithm 1:** EmbeddedCurve( $q, D_{\min}, D_{\max}$ )

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**Input:** prime integer  $q$ , minimum and maximum values of  $D > 0$

**Output:** A list of traces and discriminants of embedded elliptic curves for  $\mathbb{F}_q$

$\mathcal{L} \leftarrow \{\}$

**for**  $D$  from  $D_{\min}$  to  $D_{\max}$  **do**

**if**  $D$  is square-free and  $-D$  is a square modulo  $q$  **then**

$$s \leftarrow \begin{cases} \sqrt{-D} \bmod q & d \not\equiv 3 \pmod{4} \\ \frac{1+\sqrt{-D}}{2} \bmod q & d \equiv 3 \pmod{4} \end{cases}$$

    lift  $s$  in  $\mathbf{Z}$

$\pi \leftarrow a + bX$  the shortest non-zero element of the lattice  $\mathbb{Z}\langle q, X - s \rangle$

**if**  $\pi$  has norm  $q$  **then**

$$(t', y') \leftarrow \begin{cases} (2a, b) & \text{if } d \equiv 3 \pmod{4} \\ (2a + b, b) & \text{otherwise} \end{cases}$$

**if**  $r = q + 1 - t', r' = q + 1 + t'$  contain a large prime **then**

$$\mathcal{L} \leftarrow \mathcal{L} \cup \{(D, t', y')\}$$

**return**  $\mathcal{L}$

---

## Atkin-Morain, ECPP, and the CM method [AM93]

- internal step in ECPP: find an elliptic curve over  $\mathbf{Z}/n\mathbf{Z}$  of non-prime order of known factorization
- enumerate small  $D$  until a curve is found
- For each  $D$ , solve a norm equation  $n = A^2 + DB^2$  in  $\mathcal{O}_K$ ,  $K = \mathbf{Q}[\sqrt{-D}]$
- the curve trace is  $t' = 2A$ , check order
- Do not compute  $H_{-D}(X)$  each time, only when a good  $D$  is found

Estimated chance to solve the norm equation  $n = N_{K/\mathbf{Q}}(\eta)$ :  $\frac{1}{2h(K)}$

$\implies$  try many  $D$  until a solution is found.

$h(K)$  grows with  $D$ .

## Example with $D = 6673027$

$$q = q_{\text{BLS12-381}} = u^4 - u^2 + 1 \text{ where}$$

$$u = -0\text{xd}201000000010000 = -(2^{63} + 2^{62} + 2^{60} + 2^{57} + 2^{48} + 2^{16})$$

$$D = 6673027 \equiv 3 \pmod{4}, h(D) = 360$$

$$s \leftarrow \frac{1 + \sqrt{-D}}{2} \pmod{q}$$

$$\text{2-dim reduction (rows) GaussReduce} \begin{bmatrix} q & 0 \\ -s & 1 \end{bmatrix}$$

$$\begin{bmatrix} 125559217103576390750801819080760038445 & 148223899205865772742806981386395067 \\ -49458940538103050268164576706759590014 & 417560298539131963054131572411958320435 \end{bmatrix}$$

$$\text{1st row} \rightarrow (a, b) \text{ such that } a^2 + ab + \frac{D+1}{4}b^2 = q$$

$$4q = (2a + b)^2 + Db^2, \text{ embedded curve trace } t = 2a + b, y = b$$

$$\text{prime order } s = q + 1 - t = \frac{(t-2)^2 + Dy^2}{4},$$

$$\text{twist order } q + 1 + t = 3^2 \cdot 19^2 \cdot 953 \cdot p_{234}$$

$$\text{PARI-GP: } j(E) \pmod{q}, E: y^2 = x^3 - 3x + b_q,$$

$$b_q = 10908001762325402974914188089519822993112853370962247355940024813778856917972$$



## Related work: plain/hybrid cycles of curves

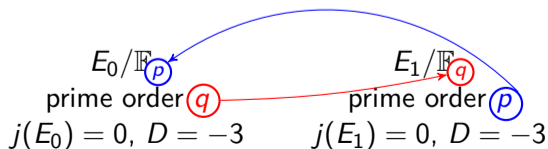
Plain cycles: 2 plain prime-order elliptic curves (no pairing)

secp256k1/secq256k1, HALO: Tweedledum/tweedledee, HALO2: Pallas-Vesta – Pasta

1. start from any prime-order elliptic curve of small-enough discriminant  $D$
2. swap scalar field order  $q$  and base field order  $p$  to get the 2nd curve parameters
3. use the CM method on  $(q, p, D)$  to get the 2nd curve coefficients

$D$  small required, SafeCurve criterion  $|D| \geq 2^{110}$  never satisfied

Hybrid cycles: a plain curve and a BN pairing-friendly curve, both prime order  
BN254-Grumpkin, BN382-plain, Pluto (BN446) - Eris. Prime-order pairing-friendly curves are very rare, no better solution than BN known



## ed\_25519 as an embedded curve

$$q = 2^{255} - 19$$

- Curve25519 in Montgomery form

$$E': y^2 = x^3 + 48662x^2 + x$$

- Ed25519 in twisted Edwards form

$$E': -x^2 + y^2 = 1 - \frac{121665}{121666}x^2y^2$$

$E'(\mathbb{F}_q)$  of order  $8r$ ,  $r$  prime

Prime  $p$ , curve  $E/\mathbb{F}_p$   
of prime order  $q$

- $D = 65012179$
- $D = 103953715$



$E'$  embedded curve of  $E$

# Outline

zk-SNARK

Elliptic Curves and Pairings

Proof-friendly curves

More embedded curves

More families of embedded curves

## Families of embedded curves with BLS12

Sanso's first work [SEH24]. Idea: design **families** of embedded curves, whose parameters are given by polynomials, like families of pairing-friendly curves. Always the same story:

1. Do it for BLS12 curves: easy doing! (fastest optimal ate pairing computation, easiest cofactor clearing, subgroup membership testing, hashing, scalar multiplication with high-dimension GLV...)
2. generalize to other pairing-friendly curves: problems arise

With previous section point of view: Solve a norm equation for  $q = x^4 - x^2 + 1$ ,  $D = 3$

Obvious solution:  $q = ((2x^2 - 1)^2 + 3(1)^2)/4$

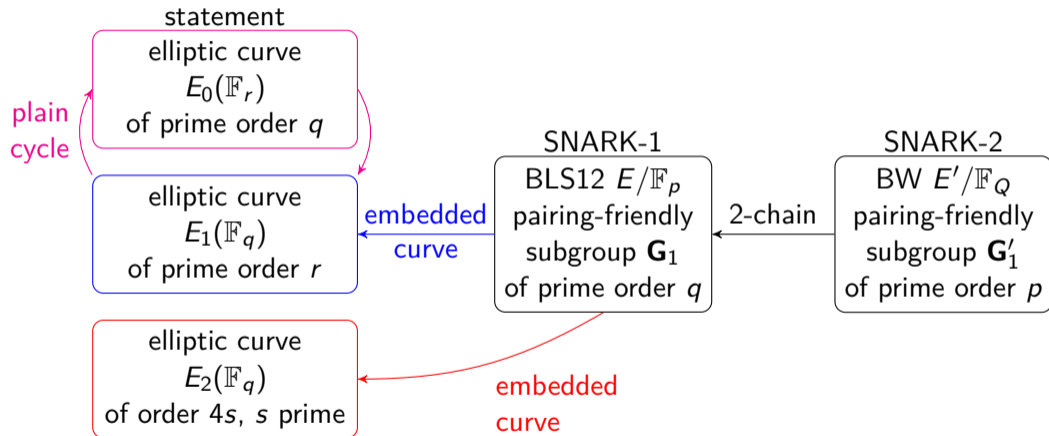
embedded curve trace:  $t' = 2x^2 - 1$ ,

embedded curve order:  $s = ((t' - 2)^2 + 3(1)^2)/4 = x^4 - 3x^2 + 3$  irreducible

*and* generating primes.

$\implies$  family!

# Generalization: Families of embedded curves for BLS12



# BLS12 with embedded curves

seed	$L$	equation $E_{\text{BLS}}/\mathbb{F}_p$	$p$ (bits)	$q$ (bits)	embedded curve equation $E_{1,2}/\mathbb{F}_q$	plain cycle curve equation $E_0/\mathbb{F}_r$
0xffff007fda000001 $2^{64} - 2^{48} + 2^{39} - 2^{29} - 2^{27} + 2^{25} + 1$	25	$y^2 = x^3 + 1$	383	256	$E_1: y^2 = x^3 + 19$ $E_2: y^2 = x^3 + 17$	$y^2 = x^3 + 7$
0xfc3ec00400000001 $2^{64} - 2^{58} + 2^{54} - 2^{48} - 2^{46} + 2^{34} + 1$	34	$y^2 = x^3 + 1$	383	256	$E_1: y^2 = x^3 + 23$ $E_2: y^2 = x^3 + 29$	$y^2 = x^3 + 29$
-0xef000ffefdfc000001 $-2^{64} + 2^{60} + 2^{56} - 2^{44} + 2^{32} + 2^{25} + 1$	25	$y^2 = x^3 + 1$	382	256	$E_1: y^2 = x^3 + 11$ $E_2: y^2 = x^3 + 17$	$y^2 = x^3 + 17$
0xdf07ffdfc000001 $2^{64} - 2^{61} - 2^{56} + 2^{51} - 2^{33} - 2^{26} + 1$	26	$y^2 = x^3 + 1$	382	256	$E_1: y^2 = x^3 + 11$ $E_2: y^2 = x^3 + 23$	$y^2 = x^3 + 7$

## Some technicalities

- $q(u) = u^4 - u^2 + 1 = \Phi_{12}(u)$  (BLS12) [SEH24]  
 $q(u) = (u^6 + 37u^3 + 343)/343$  (KSS18), with Sagemath code at [Hop20]  
 $q(u) = (u^8 + 48u^4 + 625)/61250$  (KSS16)
- Solve for  $t'(u), y'(u)$  in  $4q(u) = t'(u)^2 + Dy'(u)^2$

Solution:

- Combine Dai–Lin–Zhao–Zhou [DLZZ23] with Smith [Smi15, §4]
- BLS12 [SEH24]  $t' = 2u^2 - 1, y' = 1$
- KSS16  $t' = (31(u/5)^4 + 1)/7, y' = (-17(u/5)^4 - 1)/14$
- KSS18  $t' = -20(u/7)^3 - 1, y' = -18(u/7)^3 - 1$
- Consider the quadratic twists, 3rd and 6-th twists ( $D = 3$ ), 4-th twists ( $D = 1$ )

## Our Algorithm

$E$  has endomorphism  $\phi$ , char. poly  $\chi(X) = X^2 - t_\phi X + \deg_\phi$   
 $t_\phi^2 - 4 \deg_\phi = -Dn^2$  and  $-D$  matches  $E$ 's in  $t^2 - 4p = -Dy^2$

1.  $\lambda(x) \leftarrow$  a root of  $\chi(X) \bmod q(x)$   
e.g. if  $\chi(X) = X^2 + D$ ,  $\lambda(x) = \sqrt{-D} = (t(x) - 2)/y(x) \bmod q(x)$
2.  $U(x), V(x) \leftarrow$  half-gcd( $q(x), \lambda(x)$ )
3. with Smith's technique [Smi15, §4], reduce the matrix  
$$\begin{bmatrix} U(x) & -V(x) \\ -t_\phi U(x) + \deg_\phi V(x) & U(x) \end{bmatrix}$$
 whose determinant is  
 $\det = U^2 - t_\phi UV + \deg_\phi V^2 = \text{Res}(\chi(X), U - VX)$   
to obtain a short row  $(a_0(x), a_1(x))$
4.  $(t', y') = (a_0, a_1)$  if  $D = 1, 2 \pmod 4$ ,  
 $(t', y') = (2a_0 - a_1, a_1)$  if  $D = 3 \pmod 4$ .



## Example with KSS16

$$E_{\text{KSS16}}: y'^2 = x'^3 + ax', j = 1728, D = 1, \chi = X^2 + 1$$

1.  $q(x) = (x^8 + 48x^4 + 625)/61250$ ,  $\lambda_\phi = (x^4 + 24)/7 \pmod{q(x)}$
2.  $U, V = (1, -\lambda_\phi) = (1, -(x^4 + 24)/7)$  (no half-gcd needed)
3.  $\det \begin{bmatrix} \mathbf{b}_1 \\ \mathbf{b}_2 \end{bmatrix} = \det \begin{bmatrix} 1 & -(x^4 + 24)/7 \\ (x^4 + 24)/7 & 1 \end{bmatrix} = 1250q(x)$
4. find integers  $(i, j) \pmod{1250 = 2 \cdot 5^4}$  such that the denominator simplifies in  $(i\mathbf{b}_1 + j\mathbf{b}_2)/1250 = (i + j(x^4 + 24)/7, i(x^4 + 24)/7 - j)/1250$
5.  $x \equiv 25, 45 \pmod{70}$  by construction (KSS16)  $\implies x \equiv 5 \pmod{10} \implies 5^4 \mid x^4$ .  
Write  $x = 10x_0 + 5 = 5(2x_0 + 1) \implies$  it simplifies to  $i + 807j \equiv 0 \pmod{2 \cdot 5^4}$ .
6. enumerate over  $(i, j)$  and keep those such that  $(a_0, a_1) = (i\mathbf{b}_1 + j\mathbf{b}_2)$  satisfies  $a_0^2 + a_1^2 = q(x)$

We obtain:

$$(i, j) = (31, 17),$$

$$(t', y') = (31\mathbf{b}_1 + 17\mathbf{b}_2)/1250 = ((17(x/5)^4 + 1)/14, (31(x/5)^4 + 1)/14) .$$

## Embedded curves for KSS16

Parameters  $(t', y')$  such that  $q = (t'^2 + y'^2)/4$  with  $D = 1$ .

	$(t', y')$ s.t. $q = (t'^2 + 4y'^2)/4$	$s = q + 1 - t'$	family
$t', y'$	$(31(u/5)^4 + 1)/7, (-17(u/5)^4 - 1)/7$	$(u^8 - 386u^4 + 5^5 \cdot 17)/61250$	(yes, $c=2$ )
$-t', y'$	$(-31(u/5)^4 - 1)/7, (-17(u/5)^4 - 1)/7$	$(u^8 + 482u^4 + 5^4 \cdot 113)/61250$	(yes, $c=2$ )
$y', t'$	$(-17(u/5)^4 - 1)/7, (31(u/5)^4 + 1)/7$	$(u^8 + 286u^4 + 5^4 \cdot 113)/61250$	(yes, $c=32$ )
$-y', t'$	$(17(u/5)^4 + 1)/7, (31(u/5)^4 + 1)/7$	$(u^8 - 190u^4 + 5^5 \cdot 17)/61250$	(yes, $c=20$ )

Valid seed:  $2^{34} - 2^{32} + 2^{30} + 2^{26} - 2^5 - 2^3 - 1 = 0x343ffffd7$  (row 2), 254-bit order

## Conclusion

Inspirations from 80's and 90's papers with modern software on nowadays' CPU solve our problems!

- embedded curves as isolated points (Bandersnatch) are always possible to find with large enough  $D$
- for SNARK, additional constraint  $2^L \mid s - 1$ , larger search space  $\rightarrow$  larger  $D \rightarrow$  much longer time
- families require to change the seed  $\rightarrow$  not always possible to replace BLS12-381
- next step: combine embedded families with outer curves families like Geppetto, BW6-751?
- still unknown: embedded and *pairing-friendly* elliptic curves (only known construction: starting from the pairing-friendly curve)
- preprint at <https://inria.hal.science/hal-04750802>

Thank you for your attention.

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