

A Mathematical Characterization of the Performance of the “Multi-Slice” Projector

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Abstract

We consider an enhanced version of the well-known “Petrov-Galerkin” projection in Hilbert spaces. The proposed procedure, dubbed “multi-slice” projector, exploits the fact that the sought solution belongs to the intersection of several high-dimensional slices. This setup is for example of interest in model-order reduction where this type of prior may be computed off-line. In this note, we provide a mathematical characterization of the performance achievable by the multi-slice projector and compare the latter with the results holding in the Petrov-Galerkin setup. In particular, we illustrate the superiority of the multi-slice approach in certain situations.

Nous considérons une version améliorée de la projection de “Petrov-Galerkin” dans un espace de Hilbert. La procédure proposée, appelée “projecteur multi-tranches”, exploite le fait que la solution recherchée appartient à l’intersection de plusieurs tranches de hautes dimensions. Dans cette note, nous fournissons une caractérisation mathématique des performances atteignables par le projecteur “multi-tranches” et comparons les résultats obtenus à ceux existants dans le contexte des projections de Petrov-Galerkin. Nous illustrons ainsi la supériorité de l’approche multi-tranches dans certaines situations.

1 Introduction

Let \mathcal{H} be a Hilbert space with inner product $\langle \cdot, \cdot \rangle$ and induced norm $\|\cdot\|$. We focus on the following variational formulation:

$$\text{Find } \mathbf{h}^* \in \mathcal{H} \text{ such that } a(\mathbf{h}^*, \mathbf{h}) = b(\mathbf{h}) \quad \forall \mathbf{h} \in \mathcal{H}, \quad (1)$$

where $a : \mathcal{H} \times \mathcal{H} \rightarrow \mathbb{R}$ is a bilinear operator and $b : \mathcal{H} \rightarrow \mathbb{R}$ a linear operator. Problem (1) is quite common (it appears for example

in the weak formulation of elliptic partial differential equations) and has therefore been well-studied in the literature. In particular, it has a unique solution under mild conditions, see Lax-Milgram’s and Necas Theorems in [1, Theorems 2.1 and 2.2].

Unfortunately, solving (1) is generally an intractable problem. A popular alternative to compute an approximation of (1) is known as “Petrov-Galerkin” projection. Formally, this approach consists of approximating (1) by the following problem:

$$\text{Find } \hat{\mathbf{h}}_{\text{PG}} \in V_n \text{ such that } a(\hat{\mathbf{h}}_{\text{PG}}, \mathbf{h}) = b(\mathbf{h}) \quad \forall \mathbf{h} \in Z_m \quad (2)$$

where $V_n \subset \mathcal{H}$ is a linear subspace of dimension n and $Z_m \subset \mathcal{H}$ is a linear subspace of dimension $m \geq n$. Since the dimension of V_n and Z_m are finite, (2) admits a simple algebraic solution under mild conditions. In the literature of model reduction (see *e.g.*, [1]), Petrov-Galerkin approximation is at the core of the family of “projection-based” reduced models.

In this note we elaborate on an alternative projection procedure exploiting several approximation subspaces. Indeed, in the context of model-order reduction, standard strategies to evaluate a good approximation subspace V_n , *e.g.*, reduced basis [1] or proper orthogonal decomposition [2], typically generate a sequence of subspaces $\{V_k\}_{k=0}^n$ and positive scalars $\{\hat{\epsilon}_k\}_{k=0}^n$ such that

$$V_0 \subset V_1 \subset \dots \subset V_n \quad (3)$$

and

$$\text{dist}(\mathbf{h}^*, V_k) \leq \hat{\epsilon}_k, \quad k = 0 \dots n. \quad (4)$$

Clearly, (4) provides some useful information about the location of \mathbf{h}^* in \mathcal{H} since it restrains the latter to belong to the intersection of a set of low dimensional slices, *i.e.*,

$$\mathbf{h}^* \in \bigcap_{k=0}^n \mathcal{S}_k, \quad (5)$$

where

$$\mathcal{S}_k = \{\mathbf{h} : \text{dist}(\mathbf{h}, V_k) \leq \hat{\epsilon}_k\}, \quad k = 0 \dots n. \quad (6)$$

In standard Petrov-Galerkin projection (2), only V_n is used and the additional information provided by (5) is discarded. In this work, we consider a simple methodology to exploit the latter additional information into the projection process. More specifically, we focus on the

following optimization problem¹

$$\begin{aligned} \text{Find } \hat{\mathbf{h}}_{\text{MS}} \in \arg \min_{\mathbf{h} \in V_n} \sum_{j=1}^m (b(\mathbf{z}_j) - a(\mathbf{h}, \mathbf{z}_j))^2 \quad (7) \\ \text{subject to } \text{dist}(\mathbf{h}, V_k) \leq \hat{\epsilon}_k, \quad k = 0 \dots n, \end{aligned}$$

which can be seen as an extension of the standard Petrov-Galerkin approach. In particular, the constraints in (7) exploit the prior information (4) into the projection process: each constraint imposes that the solution belongs to some k -dimensional slice \mathcal{S}_k . Hence, in the sequel, we will dub this procedure as “multi-slice” projection.

The practical interest of the multi-slice approach has already been emphasized in several contributions. In [3, 4] we presented some applications of the multi-slice decoder to the problem of model-order reduction of parametric partial differential equations. In [5] and [6], the authors showed that multi-slice decoder can be of interest to enhance the performance of the “empirical interpolation method” or the simulation of Navier-Stokes equations. “Multi-slice” prior information of the form (5) has also been considered in [7] for data assimilation. However, in the latter contribution, the decoder considered by the authors differs from (7) since the solution is no longer constrained to belong to the low-dimensional subspace V_n .

In this note we provide a mathematical characterization of the performance achievable by the multi-slice decoder (7). More specifically, we derive an “instance optimality property” relating the projection error $\|\hat{\mathbf{h}}_{\text{MS}} - \mathbf{h}^*\|$ to the distance between \mathbf{h}^* and the different approximation subspaces V_k . Our result is presented in Theorem 2 in the next section.

2 Performance guarantees

One of the reasons which has ensured the success of Petrov-Galerkin projection is the existence of strong theoretical guarantees, *e.g.*, Cea’s Lemma [1, Lemma 2.2] or the Babuska’s Theorem [1, Theorem 2.3]. In this section we derive a similar result for the multi-slice decoder (7). The standard result associated to Petrov-Galerkin projection is recalled in Theorem 1 whereas our characterization of the multi-slice decoder (7) is presented in Theorem 2. We conclude this section by providing two examples in which the multi-slice projector leads to bet-

¹In this note we assume that constraints are available $\forall k \in \{1 \dots n\}$. All the derivations presented in this paper may nevertheless be easily extended to the case where constraints in (7) are only available for *some* $k \in \{1 \dots n\}$.

ter guarantees of reconstruction than the standard Petrov-Galerkin approach.

We first introduce some quantities of interest. First, we let $\{\mathbf{v}_j\}_{j=1}^n$ and $\{\mathbf{z}_j\}_{j=1}^m$ be orthonormal bases (ONBs) of the subspaces V_n and Z_m , respectively. We define $\{\mathbf{a}_j\}_{j=1}^m$ as the Riesz's representers of $\{a(\cdot, \mathbf{z}_j)\}_{j=1}^m$. We denote by $\{\sigma_j\}_{j=1}^n$ the set of singular values (sorted in their decreasing order of magnitude) of the Gram matrix

$$\mathbf{G} = [\langle \mathbf{a}_i, \mathbf{v}_j \rangle]_{i,j} \in \mathbb{R}^{m \times n}. \quad (8)$$

With these notations, the well-known Babuska's theorem (in a Hilbert space) can be formulated as follows:

Theorem 1 (Babuska's Theorem). *If $\sigma_n > 0$ then the solution of (2) is unique and satisfies*

$$\|\mathbf{h}^* - \hat{\mathbf{h}}_{\text{PG}}\| \leq \frac{\sigma_1}{\sigma_n} \text{dist}(\mathbf{h}^*, V_n). \quad (9)$$

See for example [8] for a proof of this result. Hereafter we provide a similar characterization of the performance of the multi-slice projector (7). In order to state our result we need to introduce the following quantities. We first define the short-hand notations²

$$\epsilon_k = \text{dist}(\mathbf{h}^*, V_k), \quad (10)$$

and

$$\gamma = \sup_{\mathbf{h} \in V_n^\perp, \|\mathbf{h}\|=1} \left(\sum_{j=1}^m \langle \mathbf{a}_j, \mathbf{h} \rangle^2 \right)^{\frac{1}{2}}. \quad (11)$$

Moreover, we define

$$\delta_j = \sum_{k=1}^n |x_{kj}| (\hat{\epsilon}_{k-1} + \epsilon_{k-1}), \quad (12)$$

where x_{kj} are the elements of the matrix \mathbf{X} appearing in the singular value decomposition of \mathbf{G} , that is $\mathbf{G} = \mathbf{U}\Lambda\mathbf{X}^T$, where $\mathbf{U} \in \mathbb{R}^{m \times m}$, $\mathbf{X} \in \mathbb{R}^{n \times n}$ are orthogonal matrices and $\Lambda \in \mathbb{R}^{m \times n}$ is the diagonal matrix of singular values $\{\sigma_j\}_{j=1}^n$.

Using these notations, our result reads:

² ϵ_k thus represents the true distance from \mathbf{h}^* to V_k . We note that this quantity is usually unknown to the practitioner. This is in contrast with $\hat{\epsilon}_k$ which represents the prior information available to the practitioner but is only an upper bound on ϵ_k .

Theorem 2. Let \mathbf{h}^* be a solution of (1) verifying (5). Then any solution $\hat{\mathbf{h}}_{\text{MS}}$ of (7) verifies

$$\|\mathbf{h}^* - \hat{\mathbf{h}}_{\text{MS}}\| \leq \begin{cases} \left(\sum_{j=\ell+1}^n \delta_j^2 + \rho \delta_\ell^2 + \epsilon_n^2 \right)^{\frac{1}{2}} & \text{if } \sum_{j=1}^n \sigma_j^2 \delta_j^2 \geq 4\gamma^2 \epsilon_n^2, \\ \left(\sum_{j=1}^n \delta_j^2 + \epsilon_n^2 \right)^{\frac{1}{2}} & \text{otherwise,} \end{cases} \quad (13)$$

where ℓ is the largest integer such that

$$\sum_{j=\ell}^n \sigma_j^2 \delta_j^2 \geq 4\gamma^2 \epsilon_n^2, \quad (14)$$

and $\rho \in [0, 1]$ is defined as

$$\rho \sigma_\ell^2 \delta_\ell^2 + \sum_{j=\ell+1}^n \sigma_j^2 \delta_j^2 = 4\gamma^2 \epsilon_n^2. \quad (15)$$

Moreover, if $\sigma_n > 0$, (7) admits a unique solution.

A proof of Theorem 2 is detailed in Section 3.

We conclude this section by particularizing the results stated in Theorems 1 and 2 to different setups. In particular, we emphasize two situations³ where the multi-slice projection has much better reconstruction guarantees than its Petrov-Galerkin counterpart. In order to ease the comparison between the bounds stated in Theorems 1 and 2, we consider the case where $\{\mathbf{a}_j\}_{j=1}^m$ is an ONB. We note that in such a case, we have $\sigma_1 \leq 1$ and $\gamma \leq 1$.

Example 1. We first assume that $\mathbf{X} = \mathbf{I}_n$ in the singular-value decomposition of \mathbf{G} . We set $\hat{\epsilon}_j = \epsilon_j$ and assume that

$$\epsilon_j = \begin{cases} 1 & j = 0 \dots n-3, \\ \epsilon^{\frac{1}{2}} & j = n-2, n-1, \\ \epsilon & j = n, \end{cases} \quad (16)$$

for some $\epsilon \ll 1$. Moreover, we let

$$\sigma_j = \begin{cases} 1 & j = 1 \dots n-3, \\ \epsilon^{\frac{1}{2}} & j = n-2, n-1, \\ \epsilon & j = n. \end{cases} \quad (17)$$

³The two setups considered below correspond to those exposed in [7, Section 3.2].

In this setup, the upper bound (9) of Theorem 1 becomes:

$$\left\| \hat{\mathbf{h}}_{\text{PG}} - \mathbf{h}^* \right\| \leq \sigma_n^{-1} \text{dist}(\mathbf{h}^*, V_n) = \epsilon^{-1} \epsilon = 1. \quad (18)$$

On the other hand, because $\mathbf{X} = \mathbf{I}$, we have

$$\delta_j = \hat{\epsilon}_{j-1} + \epsilon_{j-1} = 2\epsilon_{j-1}. \quad (19)$$

The index ℓ appearing in Theorem 2 is smaller or equal to $n-1$ since

$$\begin{aligned} \sigma_n^2 \delta_n^2 &= \sigma_n^2 (2\epsilon_{n-1})^2 = 4\epsilon^3 \ll 4\epsilon^2, \\ \sigma_{n-1}^2 \delta_{n-1}^2 &= \sigma_{n-1}^2 (2\epsilon_{n-2})^2 = 4\epsilon^2, \end{aligned}$$

and thus

$$\sigma_{n-1}^2 \delta_{n-1}^2 + \sigma_n^2 \delta_n^2 \geq 4\epsilon^2 \geq 4\gamma^2 \epsilon^2 \quad (20)$$

since $\gamma \leq 1$. The upper bound in Theorem 2 becomes

$$\begin{aligned} \left\| \mathbf{h}^* - \hat{\mathbf{h}}_{\text{MS}} \right\| &\leq (\delta_{n-1}^2 + \delta_n^2 + \epsilon_n^2)^{\frac{1}{2}}, \\ &= (4\epsilon + 4\epsilon + \epsilon^2)^{\frac{1}{2}}, \\ &\leq 3\epsilon^{\frac{1}{2}}. \end{aligned} \quad (21)$$

Hence the bound in the multi-slice setup (21) can be arbitrarily small as compared to (18) when $\epsilon \rightarrow 0$.

Example 2. We now consider $\mathbf{X} = n^{-\frac{1}{2}} \mathbf{1}_{n \times n}$ where $\mathbf{1}_{n \times n}$ is an $n \times n$ matrix of 1's. We set $\hat{\epsilon}_j = \epsilon_j$ and assume that

$$\epsilon_j = \begin{cases} \frac{1}{2} & j = 0, \\ \frac{1}{2(n-1)} & j = 1 \dots n-1, \\ \epsilon & j = n, \end{cases} \quad (22)$$

for some $\epsilon \ll n^{-1}$ (Note that we must have: $\epsilon \leq \frac{1}{2(n-1)}$ by definition). Moreover, we let

$$\sigma_j = \begin{cases} \sigma & j = 1 \dots n-1, \\ \epsilon^2 & j = n, \end{cases} \quad (23)$$

for some $1 \geq \sigma > \epsilon$ whose value will be specified below.

With these choices, the upper bound (9) of Theorem 1 becomes:

$$\left\| \hat{\mathbf{h}}_{\text{PG}} - \mathbf{h}^* \right\| \leq \sigma_n^{-1} \text{dist}(\mathbf{h}^*, V_n) = \epsilon^{-2} \epsilon = \epsilon^{-1}. \quad (24)$$

On the other hand, we have

$$\begin{aligned}
\delta_j &= \sum_{k=1}^n |x_{kj}|(\hat{\epsilon}_{k-1} + \epsilon_{k-1}), \\
&= 2n^{-\frac{1}{2}} \sum_{k=1}^n \epsilon_{k-1}, \\
&= 2n^{-\frac{1}{2}}.
\end{aligned} \tag{25}$$

By choosing σ such that (we remind the reader that $\sigma_{n-1} = \sigma$ by definition (23))

$$\sigma_{n-1}^2 \delta_{n-1}^2 + \sigma_n^2 \delta_n^2 = 4\epsilon^2, \tag{26}$$

we obtain that index ℓ appearing in Theorem 2 is smaller or equal to $n-1$ since $\gamma \leq 1$. The upper bound in Theorem 2 then reads

$$\begin{aligned}
\|\mathbf{h}^* - \hat{\mathbf{h}}_{\text{MS}}\| &\leq (\delta_{n-1}^2 + \delta_n^2 + \epsilon_n^2)^{\frac{1}{2}}, \\
&= (4n^{-1} + 4n^{-1} + \epsilon^2)^{\frac{1}{2}}, \\
&\leq 3n^{-\frac{1}{2}},
\end{aligned} \tag{27}$$

where the last inequality follows from our initial assumption $\epsilon \ll n^{-1}$.

3 Proof of Theorem 2

In this section, we provide a proof of the result stated in Theorem 2. We first note that problem (7) is equivalent to finding the minimum of a quadratic function over a closed bounded subset of V_n . A minimizer thus always exists. Moreover, the unicity of the minimizer stated at the end of Theorem 2 follows from the strict convexity of the cost function when $\sigma_n > 0$.

In the rest of this section, we thus mainly focus on the derivation of the upper bound (13). Our proof is based on the following steps. First, since $\hat{\mathbf{h}}_{\text{MS}} \in V_n$, we have that

$$\begin{aligned}
\|\mathbf{h}^* - \hat{\mathbf{h}}_{\text{MS}}\|^2 &= \|P_{V_n}(\mathbf{h}^*) - \hat{\mathbf{h}}_{\text{MS}}\|^2 + \|P_{V_n}^\perp(\mathbf{h}^*)\|^2, \\
&= \|P_{V_n}(\mathbf{h}^*) - \hat{\mathbf{h}}_{\text{MS}}\|^2 + \epsilon_n^2,
\end{aligned} \tag{28}$$

where $P_{V_n}(\cdot)$ (resp. $P_{V_n}^\perp(\cdot)$) denotes the orthogonal projector onto V_n (resp. V_n^\perp). We then derive an upper bound on $\|P_{V_n}(\mathbf{h}^*) - \hat{\mathbf{h}}_{\text{MS}}\|^2$ as follows:

- We identify a set \mathcal{D} such that $P_{V_n}(\mathbf{h}^*) - \hat{\mathbf{h}}_{\text{MS}} \in \mathcal{D}$ in Section 3.1. We then have $\|P_{V_n}(\mathbf{h}^*) - \hat{\mathbf{h}}_{\text{MS}}\|^2 \leq \sup_{\mathbf{d} \in \mathcal{D}} \|\mathbf{d}\|^2$.
- We derive the analytical expression of $\sup_{\mathbf{d} \in \mathcal{D}} \|\mathbf{d}\|^2$ as a function of the parameters $\{\epsilon_k\}_{k=1}^n$, $\{\hat{\epsilon}_k\}_{k=1}^n$ and $\{\sigma_k\}_{k=1}^n$.

Combining these results, we obtain (13)-(15).

3.1 Definition of \mathcal{D}

We express \mathcal{D} as the intersection of two sets \mathcal{D}_1 and \mathcal{D}_2 that we define in Sections 3.1.2 and 3.1.3 respectively. In order to properly define these quantities, we introduce some particular ONBs for V_n and $W_m = \text{span}\left(\{\mathbf{a}_j\}_{j=1}^m\right)$ in Section 3.1.1.

3.1.1 Some particular bases for V_n and W_m

Let

$$\mathbf{G} = \mathbf{U}\mathbf{\Lambda}\mathbf{X}^T \quad (29)$$

be the singular value decomposition of the Gram matrix defined in (8), where $\mathbf{U} \in \mathbb{R}^{m \times m}$ and $\mathbf{X} \in \mathbb{R}^{n \times n}$ are orthonormal matrices and $\mathbf{\Lambda} \in \mathbb{R}^{m \times n}$ is the diagonal matrix of singular values. We denote by $\{\sigma_j\}_{j=1}^n$ the set of singular values of \mathbf{G} sorted in their decreasing order of magnitude.

We define the following bases for V_n and W_m :

$$\mathbf{v}_j^* = \sum_{i=1}^n x_{ij} \mathbf{v}_i, \quad (30)$$

$$\mathbf{a}_j^* = \sum_{i=1}^m u_{ij} \mathbf{a}_i, \quad (31)$$

where $\mathbf{U} \in \mathbb{R}^{m \times m}$ and $\mathbf{X} \in \mathbb{R}^{n \times n}$ are the orthonormal matrices appearing in (29). We note that $\{\mathbf{v}_j^*\}_{j=1}^n$ is an ONB whereas $\{\mathbf{a}_j^*\}_{j=1}^m$ is not necessarily orthonormal. By definition, $\{\mathbf{v}_j^*\}_{j=1}^n$ and $\{\mathbf{a}_j^*\}_{j=1}^m$ enjoy the following desirable property:

$$\langle \mathbf{a}_i^*, \mathbf{v}_j^* \rangle = \begin{cases} \sigma_j & \text{if } i = j \\ 0 & \text{otherwise.} \end{cases} \quad (32)$$

3.1.2 Definition of \mathcal{D}_1

Let us define \mathcal{D}_1 as

$$\mathcal{D}_1 = \left\{ \mathbf{d} = \sum_{j=1}^n \beta_j \mathbf{v}_j^* : \sum_{j=1}^n \sigma_j^2 \beta_j^2 \leq 4\gamma^2 \epsilon_n^2 \right\}, \quad (33)$$

where γ is defined in (11). We show hereafter that $P_{V_n}(\mathbf{h}^*) - \hat{\mathbf{h}}_{\text{MS}} \in \mathcal{D}_1$.

Let us first consider the intermediate set

$$\mathcal{S} = \{\mathbf{h} : f(\mathbf{h}) \leq \gamma^2 \epsilon_n^2\}, \quad (34)$$

where $f(\mathbf{h}) = \sum_{j=1}^m (b(\mathbf{z}_j) - a(\mathbf{h}, \mathbf{z}_j))^2$ is the cost function appearing in the variational formulation of multi-slice projector (7).

Clearly $P_{V_n}(\mathbf{h}^*) \in \mathcal{S}$ because

$$\begin{aligned} f(P_{V_n}(\mathbf{h}^*)) &= \sum_{j=1}^m (b(\mathbf{z}_j) - a(P_{V_n}(\mathbf{h}^*), \mathbf{z}_j))^2 \\ &= \sum_{j=1}^m (\langle \mathbf{a}_j, \mathbf{h}^* \rangle - \langle \mathbf{a}_j, P_{V_n}(\mathbf{h}^*) \rangle)^2 \\ &= \sum_{j=1}^m (\langle \mathbf{a}_j, P_{V_n}^\perp(\mathbf{h}^*) \rangle)^2 \\ &\leq \gamma^2 \|P_{V_n}^\perp(\mathbf{h}^*)\|^2 \\ &\leq \gamma^2 \epsilon_n^2. \end{aligned} \quad (35)$$

Moreover, $\hat{\mathbf{h}}_{\text{MS}} \in \mathcal{S}$. This can be seen from the following arguments. First, $P_{V_n}(\mathbf{h}^*)$ is a feasible point for problem (7), that is

$$\text{dist}(P_{V_n}(\mathbf{h}^*), V_k) \leq \hat{\epsilon}_k \text{ for } k = 0 \dots n. \quad (36)$$

Indeed, rewriting \mathbf{h}^* as

$$\mathbf{h}^* = \sum_{j=1}^n \langle \mathbf{v}_j, \mathbf{h}^* \rangle \mathbf{v}_j + \mathbf{z}, \quad (37)$$

where $\mathbf{z} \in V_n^\perp$, we have

$$\begin{aligned}
\hat{\epsilon}_k &\geq \text{dist}(\mathbf{h}^*, V_k) \\
&= \|P_{V_k}^\perp(\mathbf{h}^*)\| \\
&= \left\| \sum_{j=k+1}^n \langle \mathbf{v}_j, \mathbf{h}^* \rangle \mathbf{v}_j + \mathbf{z} \right\| \\
&= \sqrt{\left\| \sum_{j=k+1}^n \langle \mathbf{v}_j, \mathbf{h}^* \rangle \mathbf{v}_j \right\|^2 + \|\mathbf{z}\|^2} \\
&\geq \left\| \sum_{j=k+1}^n \langle \mathbf{v}_j, \mathbf{h}^* \rangle \mathbf{v}_j \right\| \\
&= \|P_{V_k}^\perp(P_{V_n}(\mathbf{h}^*))\| \\
&= \text{dist}(P_{V_n}(\mathbf{h}^*), V_k). \tag{38}
\end{aligned}$$

The first inequality follows from our initial assumption $\mathbf{h}^* \in \cap_{k=0}^n \mathcal{S}_k$. The third equality is true because $\mathbf{z} \in V_n^\perp$. Now, since $\hat{\mathbf{h}}_{\text{MS}}$ is a minimizer of $f(\mathbf{h})$ over the set of feasible points, we have $f(\hat{\mathbf{h}}_{\text{MS}}) \leq f(P_{V_n}(\mathbf{h}^*)) \leq \gamma^2 \epsilon_n^2$ and therefore $\hat{\mathbf{h}}_{\text{MS}} \in \mathcal{S}$.

We finally show that $\hat{\mathbf{h}}_{\text{MS}} \in \mathcal{S}$ and $P_{V_n}(\mathbf{h}^*) \in \mathcal{S}$ implies $P_{V_n}(\mathbf{h}^*) - \hat{\mathbf{h}}_{\text{MS}} \in \mathcal{D}_1$. Let us first note that, if $\mathbf{h} \in V_n$, the cost function $f(\mathbf{h})$ can be rewritten as:

$$\begin{aligned}
f(\mathbf{h}) &= \sum_{j=1}^m (b(z_j) - a(\mathbf{h}, z_j))^2 \\
&= \sum_{j=1}^m (\langle \mathbf{a}_j, \mathbf{h}^* \rangle - \langle \mathbf{a}_j, \mathbf{h} \rangle)^2, \\
&= \sum_{j=1}^m (\langle \mathbf{a}_j^*, \mathbf{h}^* \rangle - \langle \mathbf{a}_j^*, \mathbf{h} \rangle)^2, \\
&= \sum_{j=1}^n (\langle \mathbf{a}_j^*, \mathbf{h}^* \rangle - \sigma_j \langle \mathbf{v}_j^*, \mathbf{h} \rangle)^2 + \sum_{j=n+1}^m \langle \mathbf{a}_j^*, \mathbf{h}^* \rangle^2, \tag{39}
\end{aligned}$$

where the third equality follows from the fact that $\{\mathbf{a}_j\}_{j=1}^m$ and $\{\mathbf{a}_j^*\}_{j=1}^m$ differ up to an orthonormal transformation; the last equality is a consequence of (32) and the fact that $\mathbf{h} \in V_n$ by hypothesis.

We note that $P_{V_n}(\mathbf{h}^*) - \hat{\mathbf{h}}_{\text{MS}}$ can be written as $\sum_{j=1}^n \beta_j \mathbf{v}_j^*$ by setting

$\beta_j = \langle \mathbf{v}_j^*, P_{V_n}(\mathbf{h}^*) \rangle - \langle \mathbf{v}_j^*, \hat{\mathbf{h}}_{\text{MS}} \rangle$. Therefore, we have

$$\begin{aligned}
\sum_{j=1}^n \sigma_j^2 \beta_j^2 &= \sum_{j=1}^n \left(\sigma_j \langle \mathbf{v}_j^*, P_{V_n}(\mathbf{h}^*) \rangle - \sigma_j \langle \mathbf{v}_j^*, \hat{\mathbf{h}}_{\text{MS}} \rangle \right)^2, \\
&= \sum_{j=1}^n \left(\sigma_j \langle \mathbf{v}_j^*, P_{V_n}(\mathbf{h}^*) \rangle - \langle \mathbf{a}_j^*, \mathbf{h}^* \rangle - \sigma_j \langle \mathbf{v}_j^*, \hat{\mathbf{h}}_{\text{MS}} \rangle + \langle \mathbf{a}_j^*, \mathbf{h}^* \rangle \right)^2, \\
&\leq 2 \sum_{j=1}^n \left(\sigma_j \langle \mathbf{v}_j^*, P_{V_n}(\mathbf{h}^*) \rangle - \langle \mathbf{a}_j^*, \mathbf{h}^* \rangle \right)^2 + 2 \sum_{j=1}^n \left(\sigma_j \langle \mathbf{v}_j^*, \hat{\mathbf{h}}_{\text{MS}} \rangle - \langle \mathbf{a}_j^*, \mathbf{h}^* \rangle \right)^2, \\
&\leq 2f(P_{V_n}(\mathbf{h}^*)) + 2f(\hat{\mathbf{h}}_{\text{MS}}), \\
&\leq 4\gamma^2 \epsilon_n^2,
\end{aligned}$$

where the first inequality follows from the standard inequality $(a+b)^2 \leq 2(a^2 + b^2)$, the second from (39), and the last one from the fact that $\hat{\mathbf{h}}_{\text{MS}} \in \mathcal{S}$ and $P_{V_n}(\mathbf{h}^*) \in \mathcal{S}$.

3.1.3 Definition of \mathcal{D}_2

Let

$$\delta_j = \eta_j + \hat{\eta}_j, \quad (40)$$

where

$$\begin{aligned}
\eta_j &= \sum_{i=1}^n |x_{ij}| \epsilon_{i-1}, \\
\hat{\eta}_j &= \sum_{i=1}^n |x_{ij}| \hat{\epsilon}_{i-1},
\end{aligned} \quad (41)$$

and the x_{ij} 's are the elements of the matrix \mathbf{X} appearing in the SVD decomposition (29). We define \mathcal{D}_2 as

$$\mathcal{D}_2 = \left\{ \mathbf{d} = \sum_{j=1}^n \beta_j \mathbf{v}_j^* : |\beta_j| \leq \eta_j \right\}. \quad (42)$$

We show hereafter that $P_{V_n}(\mathbf{h}^*) - \hat{\mathbf{h}}_{\text{MS}} \in \mathcal{D}_2$.

We first note that if \mathbf{h} is feasible for problem (7), we must have

$$|\langle \mathbf{v}_j^*, \mathbf{h} \rangle| \leq \hat{\eta}_j. \quad (43)$$

Indeed, if \mathbf{h} is feasible, the constraint $\text{dist}(\mathbf{h}, V_k) \leq \hat{\epsilon}_k$ simply writes as

$$\sum_{j=k+1}^n \langle \mathbf{v}_j, \mathbf{h} \rangle^2 \leq \hat{\epsilon}_k^2.$$

In particular, this implies that

$$|\langle \mathbf{v}_{k+1}, \mathbf{h} \rangle| \leq \hat{\epsilon}_k.$$

Using the fact that

$$\mathbf{v}_j^* = \sum_{k=1}^n x_{kj} \mathbf{v}_k,$$

we obtain (43). In a similar way, we can find that

$$|\langle \mathbf{v}_j^*, P_{V_n}(\mathbf{h}^*) \rangle| \leq \eta_j, \quad (44)$$

by using the fact that $\text{dist}(P_{V_n}(\mathbf{h}^*), V_k) \leq \epsilon_k$ from (38).

Let us now show that $P_{V_n}(\mathbf{h}^*) - \hat{\mathbf{h}}_{\text{MS}} \in \mathcal{D}_2$. We first note that $P_{V_n}(\mathbf{h}^*) - \hat{\mathbf{h}}_{\text{MS}}$ can be written as $\sum_{j=1}^n \beta_j \mathbf{v}_j^*$ by setting $\beta_j = \langle \mathbf{v}_j^*, P_{V_n}(\mathbf{h}^*) \rangle - \langle \mathbf{v}_j^*, \hat{\mathbf{h}}_{\text{MS}} \rangle$. This leads to

$$\begin{aligned} |\beta_j| &= \left| \langle \mathbf{v}_j^*, P_{V_n}(\mathbf{h}^*) \rangle - \langle \mathbf{v}_j^*, \hat{\mathbf{h}}_{\text{MS}} \rangle \right|, \\ &\leq \left| \langle \mathbf{v}_j^*, P_{V_n}(\mathbf{h}^*) \rangle \right| + \left| \langle \mathbf{v}_j^*, \hat{\mathbf{h}}_{\text{MS}} \rangle \right|, \\ &\leq \hat{\eta}_j + \eta_j = \delta_j, \end{aligned}$$

where the last inequality follows from (43) and (44).

3.2 Expression of $\sup_{\mathbf{d} \in \mathcal{D}} \|\mathbf{d}\|^2$

We consider the following problem:

$$\sup_{\mathbf{d} \in \mathcal{D}} \|\mathbf{d}\|^2 = \sup_{\boldsymbol{\beta}} \|\boldsymbol{\beta}\|^2 \quad \text{subject to} \quad \begin{cases} \sum_{j=1}^n \sigma_j^2 \beta_j^2 \leq 4\gamma^2 \epsilon_n^2 \\ |\beta_j| \leq \delta_j \end{cases}. \quad (45)$$

If $\sum_{j=1}^n \sigma_j^2 \delta_j^2 \leq 4\gamma^2 \epsilon_n^2$, the first constraint in (45) is always inactive and the solution simply reads

$$\sup_{\mathbf{d} \in \mathcal{D}} \|\mathbf{d}\|^2 = \sum_{j=1}^n \delta_j^2. \quad (46)$$

If $\sum_{j=1}^n \sigma_j^2 \delta_j^2 \geq 4\gamma^2 \epsilon_n^2$, the solution of (45) is given by

$$\sup_{\mathbf{d} \in \mathcal{D}} \|\mathbf{d}\|^2 = \sum_{j=\ell+1}^n \delta_j^2 + \rho \delta_\ell^2, \quad (47)$$

where ℓ is the largest integer such that

$$\sum_{j=\ell}^n \sigma_j^2 \delta_j^2 \geq 4\gamma^2 \epsilon_n^2, \quad (48)$$

and $\rho \in [0, 1]$ is defined as

$$\rho \sigma_\ell^2 \delta_\ell^2 + \sum_{j=\ell+1}^n \sigma_j^2 \delta_j^2 = 4\gamma^2 \epsilon_n^2. \quad (49)$$

This can be seen by verifying the optimality condition of problem (45). We note that problem (45) is the same (up to some constants) to the one considered in [7, Section 3.1]. The solution (47) is therefore similar, up to some different constants, to the one obtained in that paper.

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